

Simulation course FYSM350

Fortran : Vesa Apaja email: vesa.apaja@gmail.com

MD : Hannu Häkkinen

MC : Juha Merikoski

www info:

Course homepage:

<http://www.phys.jyu.fi/homepages/merikosk/Simu2006.html>

(a link to this page is in korppi)

Useful material:

<http://www.csc.fi/oppaat/f95/>

Introduction

- History

“Dad's FORTRAN”

First version born late 50's at IBM

First standard 1966 : FORTRAN 66

Second standard 1977: FORTRAN 77

----- prehistory -----

Third standard 1990 : Fortran 90

Fourth standard 1995 : Fortran 95

Fifth standard 2003 : Fortran 2003

Describing his early work on FORTRAN, John Backus said:-

We did not know what we wanted and how to do it. It just sort of grew. The first struggle was over what the language would look like. Then how to parse expressions - it was a big problem and what we did looks astonishingly clumsy now....

Free compilers

Intel Fortran (ifort or ifc)

<http://www.intel.com/software/products/compilers/flin/noncom.htm>

- Windows or linux, any Intel or AMD processor
- Optimising f95 compiler with 2003 “readiness”
- Free for non--commercial use; Windows version is **only evaluation copy**
- After registering you get a license number
- Easy to install (linux: less than 5 mins)

Gnu g95

<http://g95.sourceforge.net/>

Source code available - written in C :^)

precompiled binaries for

Linux x86, Cygwin x86, Windows x86, Powerpc, FreeBSD x86,
Sparc Solaris, Linux IA64, Linux x86_64/EMT64, Linux Alpha,
Irix Mips *etc.*

Commercial compilers

Often linked to commands **f90** or **f95**

- Pathscale **pathf90**
- PGI (Portland group) **pgf90**
- MIPSPRO (SGI Irix machines)
- Lahey (Fujitsu) **lf95**
- Absoft
- Compaq Fortran (Digital Fortran)

Intel: Compiler options

Try **ifort -help** or **man ifort**

Development phase:

ifort -check all program.f90

Production phase:

Always try : **-fast**

ifort -O3 program.f90

Use always at least this

ifort -tpp7 -xP

pentium 4 processor

Expect output lines “*remark: LOOP WAS VECTORIZED*”

g95: Compiler options

Development phase:

g95 -Wall -std=f95 program.f90

Production phase:

g95 -O3 -std=f95 program.f90

Try also : **-O2 , -funroll-all-loops -mtune=pentium4**

(practically all options of the gcc compiler apply)

First program

```
program test
```

```
implicit none
```

```
integer:: i
```

```
real (kind(1.d0)) :: x,y
```

```
complex (kind(1.d0)) :: c
```

```
x = 5.d0
```

```
y = 20.d0
```

```
c = cmplx(x,y,kind(1.d0))
```

```
print*,x,y,c
```

```
end program test
```

Good habit: use this always if your program is meant for **serious use**

Unused variable: waste of space

Double precision:
~ 16 decimal accuracy

Internal function “**cmplx**”

Convert two double precision numbers to a complex number

output:

```
5.0000000000000000    20.0000000000000000  
(5.0000000000000000,20.0000000000000000)
```

Formatted output

Most important format codes: examples

2f15.5 two floating point numbers, use 15 characters and give 5 decimals

f0.10 one floating point numbers, 10 decimals, as much field as it takes

a20 one 20 characters long field

d15.10 one number of the exponent type **10.d12** **5x** 5 spaces

Bad format codes may

- cause the program to terminate
- clobber the output to a useless form:

```
write(*,'(3f0.3)') 0.1234d0,0.5678d0,0.7890d0      0.1230.5680.789
```

Examples:

```
write(*,'(d15.5)') 1.123456789d12
```

0.11235D+13

```
write(*,'(g15.5)') 1.123456789d12
```

0.11235E+13

Field is

```
write(*,'(f15.5)') 1.123456789d12
```

← too

```
write(*,'(e15.5)') 1.123456789d12
```

0.11235E+13

short

Simple formatted output to a file

```
program test
```

```
implicit none
```

```
integer, parameter:: dbl=kind(1.d0)
```

```
integer:: i
```

```
real (dbl), parameter:: dx=0.1d0
```

```
real (dbl) :: x
```

```
write(12,'(2a15)') 'x','sin(x)
```

```
do i = 1, 10
```

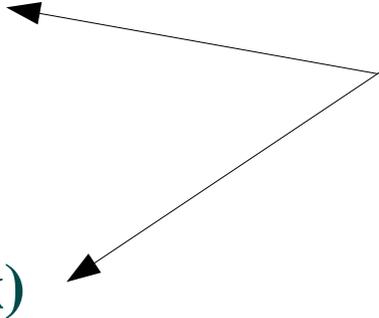
```
  x = (i-1)*dx
```

```
  write(12,'(2f15.10)') x,sin(x)
```

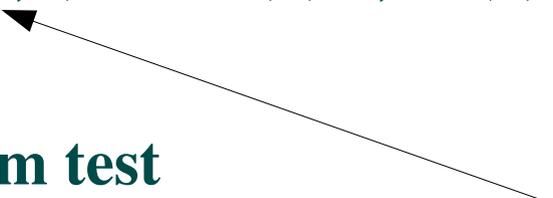
```
end do
```

```
end program test
```

**“Automatic” file I/O:
Output to file fort.12**



Unit number



Unit * : default
Unit 5 : keyboard
Unit 6 : screen

Formatted output cont'd

- Format strings in write statements:

```
write(55,"result = ",f15.10)  res
```

- For frequently occurring formats use a separate **format** statement:

...

```
kin_E = 12.66666666666666d0
```

```
pot_E = -1.44444444444444d0
```

```
write(*,900) ' kinetic energy',kin_E
```

```
kinetic energy = 12.66666666667
```

```
write(*,900) 'potential energy',pot_E
```

```
potential energy = -1.44444444444
```

...

```
900 format(1x,a20," = ",f15.10)
```

...

Basic loop structures

```
j = 6
do i = 1, 10
  print*,i
  if(i==j) exit
end do
```

```
do i = 1, 10
  if(i<5) cycle
  print*,i
end do
```

```
i = 0
do while (i<3)
  i = i + 1
end do
```

very useful !

e.g. In force calculations particles don't exert force on themselves, so for *i*:th particle you need to skip particle *j* if *j=i*

If something goes wrong use **stop**:

```
subroutine mysub(x)
...
if(x<0.d0) stop 'mysub: negative x'
...
```

In bigger programs it's a good idea to let statements tell where they are ...

... and what exactly went wrong

Formatted output on screen

```
program test
```

```
implicit none
```

```
integer, parameter:: dbl=kind(1.d0)
```

```
integer:: i
```

```
real (dbl), parameter:: dx=0.1d0
```

```
real (dbl) :: x
```

```
write(*,'(2a15)') 'x','sin(x)'
```

```
do i = 1, 10
```

```
  x = (i-1)*dx
```

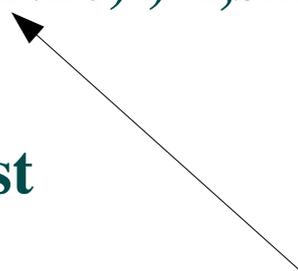
```
  write(*,'(2f15.10)') x,sin(x)
```

```
end do
```

```
end program test
```

	x	sin(x)
	0.0000000000	0.0000000000
	0.1000000000	0.0998334166
	0.2000000000	0.1986693308
	0.3000000000	0.2955202067
	0.4000000000	0.3894183423
	0.5000000000	0.4794255386
	0.6000000000	0.5646424734
	0.7000000000	0.6442176872
	0.8000000000	0.7173560909
	0.9000000000	0.7833269096

Output format



Subroutines and functions

```
subroutine swap(a,b)
```

```
  implicit none
```

```
  real(kind(1.d0)):: a,b,tmp
```

```
  tmp = a
```

```
    a = b
```

```
    b = tmp
```

```
end subroutine swap
```

in calling program:

```
...
```

```
call swap(c,d)
```

```
...
```

```
real(kind(1.d0)) function poly(x)
```

```
  implicit none
```

```
  real(kind(1.d0)) :: x
```

```
  poly = x**6+x**3-x**2+x+5.d0
```

```
end function poly
```

in calling program:

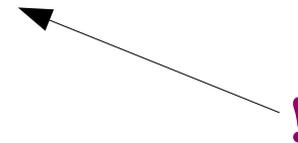
```
...
```

```
integer, parameter:: dbl=kind(1.d0)
```

```
real(dbl),external:: poly
```

```
print*,poly(2)
```

```
...
```



- **Arguments: “Pass by reference”, the content of a variable can *really* be changed**
- **Since Fortran 90: functions can return array valued results**
example: internal function for matrix multiplication

How to make a function that returns an array

Advanced feature

```
program test
```

```
implicit none
```

```
interface
```

```
function sqroot(x) result(y)
```

```
real(kind(1.d0)):: x(:)
```

```
real(kind(1.d0)):: y(size(x))
```

```
end function sqroot
```

```
end interface
```

```
integer:: i
```

```
real(kind(1.d0)):: x(5)
```

```
do i = 1, 5
```

```
  x(i) = 5.d0*i
```

```
end do
```

```
write(*,('  x ',5f10.3)) x
```

```
write(*,('sqrt(x)',5f10.3)) sqroot(x)
```

```
end program test
```

x	5.000	10.000	15.000	20.000	25.000
sqrt(x)	2.236	3.162	3.873	4.472	5.000

```
function sqroot(x) result(y)
```

```
implicit none
```

```
integer:: i
```

```
real(kind(1.d0)) :: x(:),y(size(x))
```

```
do i = 1, size(x)
```

```
  y(i) = sqrt(x(i))
```

```
end do
```

```
end function sqroot
```

The result is *not*
called sqroot

No sqroot=...

**Interface: need to tell the compiler
that we want back an array y(:)**

One cannot define

```
real(kind(1.d0)), external:: sqroot(5)
```

Warning: don't compute sqrt like this; sqrt(x) does exactly the same !

Intrinsic functions are “elemental”

and know how to operate on many types of data

program test

implicit none

integer:: i,j

real(kind(1.d0)):: array(4,4)

do i = 1, 4

do j = 1, 4

call random_number(array(i,j))

end do

end do

write(*,'(4(4f5.2,/))') array,sin(array)

end program test

4x4 array

0.00 0.96 0.80 0.09

0.03 0.84 0.83 0.89

0.35 0.34 0.35 0.70

0.67 0.92 0.87 0.73

Original
array

built-in random number generator

0.00 0.82 0.71 0.09

0.03 0.74 0.74 0.78

0.35 0.33 0.34 0.64

0.62 0.79 0.77 0.67

sin() of every
element
taken separately

/ means newline

print*,exp(1.0)

2.718282

print*,exp(1.d0)

2.71828182845905

print*,exp(cmplx(1.d0,2.d0))

(-1.131204,2.471727)

compilers don't like this: exp(1)

**Functions are mostly generic:
Argument type determines
the output type
Exception: cmplx() !**

Naming loops

```
iloop: do i = 1, 3
  jloop: do j = 1, 3
    kloop: do k = 1,3
      print*,i,j,k
      if(i==k) exit jloop
    end do kloop
  end do jloop
end do iloop
```

i	j	k	
1	1	1	↖ exit jloop
2	1	1	
2	1	2	↖
3	1	1	
3	1	2	
3	1	3	↖

Program can **jump out of several nested loops**,
not just the present loop (**if(i==k) exit** would do that)

Equally important : In a long program it's hard to tell what loop
does a bare **end do** really end.

Passing functions as arguments

Often you need to tell a function or a subroutine to use a certain function, which is passed to them as an argument. A typical example is an integral.

```
! -----
! tests subroutine integral
! -----
program test
  implicit none
  real(kind(1.d0)):: pi,res,step,a,b
  intrinsic:: dsin
  pi = 4.d0*atan(1.d0)
  step = 0.1d0
  a = 0.d0
  b = pi
  do
    call integral(dsin,a,b,step,res)
    write*,'(2es25.16)' step,res
    step = 0.1d0*step
    if(step<1.d-8) exit
  end do
end program test
```

```
! -----
! computes the definite integral
! of f(x) from a to b with step dx
! -----
subroutine integral(f,a,b,dx,res)
  implicit none
  real(kind(1.d0)), intent(in):: a,b,dx
  real(kind(1.d0)),intent(out):: res
  real(kind(1.d0)), external :: f
  real(kind(1.d0)):: x
  res = 0.d0
  x = a
  do
    res = res + f(x)
    x = x + dx
  if(x>b) exit
  end do
  res = dx*res
end subroutine integral
```

intent(in) : variables we are *not allowed* to change (by mistake) inside the subroutine.
intent(out) : variables we *must* give a value
external :: f
 tells the compile that **f** is not a table, but a function

1.000000000000000001E-01	1.9995479597125969E+00
1.000000000000000002E-02	1.9999900283082577E+00
1.000000000000000002E-03	1.9999999540411613E+00
1.000000000000000003E-04	1.9999999986726047E+00
1.000000000000000004E-05	1.9999999999843638E+00
1.000000000000000004E-06	1.9999999999895848E+00
1.000000000000000005E-07	2.0000000005709588E+00
1.000000000000000005E-08	2.0000000077288904E+00

As an advanced feature, the next page explains the reason for using **dsin** and not **sin**

Passing functions as arguments

Advanced feature

After just stating that functions are generic ... compilers are no infallible

program test

```
implicit none
real(kind(1.d0)):: x,sini
intrinsic:: dsin ← !
x = atan(1.d0) ! this is pi/4
print>(' argument ',f15.5)",x
print>(' sin(x) ',f15.5)",sin(x)
print>('own sin(x) ',f15.5)",sini(dsin,x)
end program test
```

```
real(kind(1.d0)) function sini(f,x)
implicit none
real(kind(1.d0)),intent(in):: x
real(kind(1.d0)),external:: f
sini = f(x)
end function sini
```

argument	0.78540
sin(x)	0.70711
own sin(x)	0.70711

Try with **sin** instead of **dsin**:

- Intel ifort version 9.0 still compiles:

argument	0.78540
sin(x)	0.70711
own sin(x)	-0.74464 rubbish!

- g95 refuses to compile with **sin**:

Internal error: g95_get_typednode(): Bad typespec

But: **real(kind(1.d0)), intrinsic:: sin**

and it *will* compile, with the same wrong result :(

Conclusion: Be careful to check that the program sends the correct intrinsic function!

The compiler may guess it wrong and send a single precision routine to a double precision routine

Modules

Modules replace the old **common** block structure: more versatile, easier to debug (both for you and the compiler). Typically one defines global parameters, data structures and utility functions in modules.

example of file **modules.f90**:

```
module kinds
  integer, parameter:: dbl=kind(1.d0)
end module kinds

module parameters
  integer, parameter:: natoms = 100  ! # of atoms
  integer, parameter:: ndim = 3      ! dimension
  integer, parameter:: ntherm=100   ! thermalization period
end module parameters
```

```
module positions
  use kinds
  use parameters, only: natoms, ndim
  real(dbl) :: r(ndim,natoms)
end module positions
```

module positions “inherits” the module **kinds** (so that we can use **dbl**)

← only attribute: use only part of the module **parameters**

outline of file **main.f90**:

```
program main
  use kinds
  use parameters, only: ntherm
  implicit none
  ...
end program main
```

Important: modules must be compiled before they can be used in a **use statement**
typically compile using order **f90 modules.f90 main.f90**

Modules may contain functions or subroutines:

keyword **contains**

For example spline subroutines and random number generators are easy to use if they are in a module

Outline of a spline module:

```
module splines
contains
  subroutine spline(x,y,n,y2)
    ...
  end subroutine spline
  subroutine splint(x,y,y2,n,xx,fx)
    ...
  end subroutine splint
end module splines
```

Outline of a random number module:

```
module random
contains
  subroutine rng(x,n)
    ...
  end subroutine rng
end module random
```

First simulation – with graphics !

File `gnusimu_example.f90`

This simulates diffusion of non-interacting particles in 2D, starting from all particles in the same point (0,0).

Usage:

in unix shell, type `mkfifo fifo` , then `a.out &` and then `gnuplot < fifo`

Features:

- **inquire statement** : used to make sure the file `fifo` exists
- **open the file 'fifo' with status='old'**: the file has to exist already
 status may be 'old' (must exist), 'new' (created now), 'unknown' (we don't know nor care if it is there or not)
- write `gnuplot` command (unset key *etc.*) to the file
- write particle positions in table `r(2,n)` to file one by one; actually `r(2,n)` is **n two-dimensional vectors**
- get a random shift `dr(2)` using the built-in random number generator `random_number` from range [0,1]
- shift `dr` by -0.5, because we want a random number [-0.5:0.5]
- move each particle amount `dr` : add **vector `dr`** to **coordinate vectors `r`**
- **periodic boundary conditions**: if a particle leaves the box, it re-enters from the other side
 realized using the **where construct**: `where(condition) do-something`

`where(r(:,i)<-1.d0) r(:,i) = r(:,i)+2.d0`

tests elements separately

applies operation to the element, which met the condition

Example:

`integer:: a(2)`

`a(1) = 60`

`a(2) = 100` **now a = 60, 100**

`where(a>70) a = a-10` **now a = 60, 90**

Limits of data types

Intrinsic functions **tiny** or **huge** tell how small or big values a data type can have.

```
program main
```

```
implicit none
```

```
integer :: i
```

```
integer(8) :: j
```

```
real:: x_single
```

```
real(kind(1.d0)):: x
```

```
print*, 'max i (integer) ', huge(i)
```

```
print*, 'max j (integer kind 8) ', huge(j)
```

```
print*, 'max x (single precision) ', huge
```

```
(x_single)
```

```
print*, 'max x (double precision) ', huge(x)
```

```
print*, 'min x (single precision) ', tiny(x_single)
```

```
print*, 'min x (double precision) ', tiny(x)
```

Very rarely needed, but it's good to know a big integer is already there

max i (integer) 2147483647

max j (integer kind 8) 9223372036854775807

max x (single precision) 3.4028235E+38

max x (double precision)

1.797693134862316E+308

min x (single precision) 1.1754944E-38

min x (double precision) 2.225073858507201E-308

```
end program main
```

This is what happens after an integer exceeds it's upper limit:

```
program main
```

```
implicit none
```

```
integer :: i
```

```
i = 1
```

```
do
```

```
if(i<0) exit
```

```
i = i + 1
```

```
end do
```

```
print*,i
```

```
end program main
```

This test should always fail, so the loop should never end!

-2147483648

positive integer turns negative !

Overflowing integers cause erratic behaviour

$2^{31} = 2147483648$, so in binary the biggest *positive* i 2147483647 is

11111111111111111111111111111111 (30 1's)

Limits of data types : cont'd

In double precision you can express numbers as small as 2.225073858507201E-308
and as big as 1.797693134862316E+308

Can there ever be a situation when these are not enough? The universe has about 10^{80} atoms !?

**BUT: Your variables may be able to express the single numbers,
but if you try to compute with *both* big and small numbers you get trouble**

```
program main
implicit none
real(kind(1.d0)):: x,y
x = 1.d23
y = 1.d-11
print*,x
print*,y
print*,(x-y)/x,y/x
end program main
```

definitely NOT too much asked for double precision

9.999999999999999E+022

9.999999999999999E-012

1.000000000000000 9.999999999999999E-035

**not quite the result,
we lost y completely**

**computer couldn't store so small *deviation* from 1
in the form $a \cdot 10^b$ with 16 decimals in a**

Conclusion: Always choose units so that numerically stored numbers are about 1

Physical constants may be huge or tiny: Planck's constant, Avogadro number, Ångström *etc.*
Avoid using their tabulated SI unit values as such.

Allocating memory: automatic and manual allocation

```
subroutine swap(A,B,n,m)
```

```
  implicit none
```

```
  integer, intent(in):: n,m
```

```
  real(kind(1.d0)):: A(n,m),B(n,m) ← Matrices to be swapped
```

```
  real(kind(1.d0)):: C(n,m) ← automatically allocated table:  
  C = A                               after leaving the subroutine C  
  A = B                               is also deallocated automatically
```

```
  B = C
```

```
end subroutine swap
```

Same operation using manual allocation:

```
subroutine swap(A,B,n,m)
```

```
  implicit none
```

```
  integer, intent(in):: n,m
```

```
  real(kind(1.d0)):: A(n,m),B(n,m)
```

```
  real(kind(1.d0)),allocatable:: C(:,:) ← Define that C will be have two indices
```

```
  allocate(C(n,m)) ← Allocate C
```

```
  C = A
```

```
  A = B
```

```
  B = C
```

```
  deallocate(C) ← This may be left out, C will be deallocated anyhow  
end subroutine swap                               when the program leaves the subroutine
```

After leaving, the content of C is forgotten

The save attribute

Usually functions and subroutines automatically deallocate internal variables and their content is lost. To have a variable that keeps its content between subsequent call to a subroutine use the **save** attribute. A typical example is a *counter*, that counts how many function calls have been made.

```
real(kind(1.d0)) function f(x)
```

```
implicit none
```

```
real(kind(1.d0)), intent(in):: x
```

```
integer, save:: count=0
```

```
f = sin(x)
```

```
count = count + 1
```

```
if(count>100) stop 'more than 100 calls'
```

```
end function f
```

**Variable count will remain in memory between calls to function f
notice how count is initialized to 0**

Another example is tables you want to allocate only once:

```
subroutine collect(vect,n)
```

```
implicit none
```

```
integer, intent(in):: n
```

```
real(kind(1.d0)), intent(in):: vect(n)
```

```
real(kind(1.d0)), allocatable, save, dimension(:):: vsum
```

```
integer, save:: count=0
```

```
if(count==0) allocate(vsum(n))
```

```
vsum = vsum + vect
```

```
count = count + 1
```

```
write(*,*) count,vsum/count
```

```
end subroutine collect
```

Reading in data : read_1_example.f90

Very simple data input

```
program test
  implicit none
  real(kind(1.d0)):: x
  print*,'give x'
  read*,x
  print*,'ok, got ',x
end program test
```

← **Waits for a number,
usually crashes on character input**

Reading in data : read_2_example.f90

```
! -----  
! Example program  
! input from a file  
! -----
```

```
program test  
  implicit none  
  integer:: io,i,k  
  integer, parameter:: n = 10  
  real(kind(1.d0)):: x(n),y(n)
```

```
!  
! read in x  
!
```

```
open(12,file='input',status='old')  
do i = 1, n  
  read(12,*,iostat=io) x(i),y(i)  
  if(io/=0) then  
    print*, 'EOF while reading data'  
    exit  
  end if  
end do  
print*, 'got ', i-1, ' value pairs x,y'  
do k = 1, i-1  
  write(*, '(2f7.3)') x(k),y(k)  
end do  
end program test
```

	If the file 'input' contains	the output is
		EOF while reading data
		got 8 value pairs x,y
	1 1.23	1.000 1.230
	2 2.22	2.000 2.220
	3 3	3.000 3.000
	4.5 4	4.500 4.000
	5.2 5	5.200 5.000
	6.1 6	6.100 6.000
	7.2 7	7.200 7.000
	8.1 8	8.100 8.000

Status is 'old': we don't want to create the data file with this program, just read it in.

Iostat=io puts to variable io a non-zero value if reading fails, in this case the file ends before all n data has been read in. The program prints out all values it got.

Extra functionality: `rewind(12)` rewinds the device 12 back to beginning, in this case the contents of the file 'input' could be read again.

Reading in parameters : namelist_example.f90

Namelist is a neat way to read in simulation parameters:

program test

implicit none

integer:: natoms=100, ndim=3

real(kind(1.d0))::

boxx=0.d0,boxy=0.d0,boxz=0.d0

namelist /para/ natoms, ndim, boxx,boxy,boxz

read(*,nml=para)

write(*,'(a21)') 'simulation parameters'

write(*,'(21('='))')

write(*,98) 'natoms',natoms

write(*,98) 'ndim',ndim

write(*,99) 'boxx',boxx

write(*,99) 'boxy',boxy

write(*,99) 'boxz',boxz

98 format(a10," = ",i8)

99 format(a10," = ",f8.3)

end program test

Default values

List of names that can appear in the namelist

Read in the namelist

Example: if the file 'data' contains

```
&para
ndim = 2,
natoms = 100,
boxx = 12.d0,
boxy = 10.d0
/
```

then the output of the command

a.out<data is

simulation parameters

=====

natoms = 100

ndim = 2

boxx = 12.000

boxy = 10.000

boxz = 0.000

The order of items is irrelevant, but the names have to match those given in the namelist statement

Reading and writing binary data files: binary_data_example.f90

A very effective way to store big chunks of data without losing accuracy in conversions

program test

implicit none

integer:: i,k

integer, parameter:: n = 10

real(kind(1.d0)):: x(n),y(n)

!

! fill table x

!

do i = 1, n

 x(i) = dble(i)

end do

!

! write x to a binary data file

!

open(12,file='input.bin',form='unformatted')

write(12) x

rewind(12)

read(12) y

write(*,*) "Binary date write/read test"

write(*,'(2a10)') 'wrote','read'

do k = 1, n

 write(*,'(2f10.3)') x(k),y(k)

end do

end program test

Creating a temporary binary file, named automatically:

```
open(12,status='scratch',form='unformatted')
```

```
write(12) myhugedataarray
```

```
rewind(12)
```

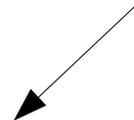
```
... DON'T CLOSE THE SCRATCH FILE YET!
```

```
read(12) anotherhugearray
```

```
close(12)
```

form = 'unformatted' means binary data

form = 'formatted' means ascii data



No format for unformatted data :^)

we just dump x to disk

Binary date write/read test

wrote	read
-------	------

1.000	1.000
-------	-------

2.000	2.000
-------	-------

3.000	3.000
-------	-------

4.000	4.000
-------	-------

5.000	5.000
-------	-------

6.000	6.000
-------	-------

7.000	7.000
-------	-------

8.000	8.000
-------	-------

9.000	9.000
-------	-------

10.000	10.000
--------	--------

Endianness: a binary data issue

This is too much for beginners, just skip it unless you already know what endianness or bit order means.

It's just a matter of consensus in whether binary data is written more significant or less significant bits first. In this respect Linux machines are usually configured to be **natively** *little endian*, but many other Unix machines are *big endian*. As a result, binary data written in one machine is unreadable (or reads in wrong) in the other unless you fix the problem. The best way to transfer binary data files among different machines is to use same endianness in both, defined via a specific *compiler option*:

Intel: compile programs using

ifort -convert big_endian program.f90

g95 : **g95 -fendian=big program.f90**

or

set the environment variable G95_ENDIAN

in bash shell : **export G95_ENDIAN=big**

in tcsh or csh : **setenv G95_ENDIAN big (or: BIG)**

then as usual **g95 program.f90**

Makefile

Helps compilation and maintenance of a program group spread over several files.

make looks for a file called **Makefile** or **makefile** - in which order seems to vary.

The command **make** takes the first target from the makefile and does what the target needs.

- **TARGET** is what we want to make
 - A binary executable
 - an object file
 - Clean up, delete old object files, temporary data files *etc.*
- **DEPENDENCY LINES** tell what the target depends on
 - If you change a module, you *MUST* recompile every routine that depends on it
- **RULES** how to create a file
 - The binary executable depends on object files, we need to tell how object files *.o are made from the program files *.f90

Syntax

TARGET: DEPENDENCIES

[press tab] **RULE**



Example:

```
main.o : main.f90
```

```
f90 -c main.f90
```

Why use a makefile: make compares the modification times (timestamps) of the targets and their dependencies and updates them *if necessary* by recompiling. Changes made to a routine will be made known to all routines that depend on it.

=> If you have 23 subroutines and you edit one of them you don't recompile every one of them.

Makefile cont'd

If you don't want to write the Makefile yourself, use a perl script **“makemake”**: It will look through all programs in the directory and create a Makefile. This is what came out for the example program:

PROG = simu

SRCS = main.f90 modules.f90 sub1.f90 sub2.f90

OBJS = main.o modules.o sub1.o sub2.o

LIBS = If you use libraries this is where to add them

CC = cc ← Unnecessary but harmless

CFLAGS = -O ←

FC = f77 ←

FFLAGS = -O

F90 = f90

F90FLAGS = -O

LDFLAGS = -s Except this -s option the Makefile works just fine.

all: \$(PROG)

\$(PROG): \$(OBJS)

\$(F90) \$(LDFLAGS) -o \$@ \$(OBJS) \$(LIBS)

clean:

rm -f \$(PROG) \$(OBJS) *.mod

.SUFFIXES: \$(SUFFIXES) .f90

.f90.o:

\$(F90) \$(F90FLAGS) -c \$<

main.o: modules.o

 Dependencies came out exactly as they should :^)

sub1.o: modules.o

User-defined data types

Single numbers, characters and tables are often insufficient to express data in a natural way..
Often the program can be made more readable (=easier to debug) with better data structures.

Syntax:

type mytype

variable definition (integers, character etc.)

end type mytype

Usage:

type(mytype):: one_of_mytype

Members are referred with the % sign,
one_of_mytype%member

Hint:

in your mind, read the % sign as if it were genitive:

Mikko%father means “Mikko's father”

atom%type means “atom's type”

atom(i)%coord(k) means “atom i's coordinate k”

```
program test
```

```
implicit none
```

```
type human
```

```
integer::nchildren
```

```
character(15) :: mother, father, child(20)
```

```
end type human
```

```
type(human):: Mikko
```

```
Mikko%mother="Sarah"
```

```
Mikko%father="Jack"
```

```
Mikko%child(1) = "Roger"
```

```
Mikko%child(2) = "Gillian"
```

```
Mikko%nchildren = 2
```

```
print*,"Mikko's father is ",Mikko%father
```

```
print*,"Mikko's children are ",Mikko%child(1:Mikko%nchildren)
```

```
end program test
```

```
Mikko's father is Jack
```

```
Mikko's children are Roger      Gillian
```

BLAS Basic Linear Algebra Subroutines

Level 1 BLAS: scalar, scalar-vector or vector-vector operations

Level 2 BLAS: vector-matrix operations

Level 3 BLAS: matrix-matrix operations

Compile yourself (www.netlib.org) or use vendor-specific implementations:

AMD [ACML](#)

Apple [Velocity Engine](#)

Compaq [CXML](#)

Cray [libsci](#)

HP [MLIB](#)

IBM [ESSL](#)

Intel [MKL](#)

NEC [PDLIB/SX](#)

SGI [SCSL](#) or [complib.sgimath](#)

SUN [Sun Performance Library](#)

Pointers

Two usages of pointers:

- Pointers can or sometimes must be used to replace ordinary allocatable variables
 - User-defined data structures cannot have the allocatable attribute, but must be defined using pointers
 - A function *cannot* return an allocatable table, but it can return
 - (i) an automatically allocated table
 - (ii) a pointer that has been dynamically allocated (example on the next page)
 - Pointers can be used to make linked lists
- Pointers as aliases to already allocated variables
 - as a name for a part of table

Example:

N particles have 3 coordinates each and there are **Nconf** sets of these **N** particle simulations, everything stored in a big table **r_all**

```
real(kind(1.d0)), allocatable, target:: r_all(:, :, :)  
real(kind(1.d0)), pointer:: r(:, :)  
allocate (r_all(Nconf, 3, N))  
r => r_all(iconf, :, :)    ! r is now the coordinates of the set iconf  
...  
all that is done here to r(k,i) affects automatically r_all(iconf,k,i)  
...
```

Pointers cont'd

This example shows how to allocate memory in a function. A table cannot have the “allocatable” attribute in a function, so one *must* use a pointer.

program test

implicit none

interface

function func(n) result(tab)

integer,intent(in):: n

real(kind(1.d0)), pointer:: tab(:)

end function func

end interface

integer,parameter:: dbl=kind(1.d0)

real(dbl), pointer:: ptr(:)

ptr => func(6)

print*,ptr

deallocate(ptr)

end program test

Things that are different in this program:

- No “external” attribute to function func
- No “real(kind(1.d0))” typing of function func (not returning a single number)
- **Explicit interface:** func deals with a pointer and that needs to be told to the compiler
- *Any* kind of call to function func allocates memory, for example just `print*,func(3)` would allocate memory. Better free the memory in the end to be sure there is no memory leakage
- Give the name `prt` to the newly made table to be able to use it later
- Function func return a pointer, that points to an allocated memory block

← Here we create a 6 element table, allocate memory for it and set all elements to 1.d0
Output: 1. 1. 1. 1. 1. 1.

function func(n) result(tab)

implicit none

integer,intent(in):: n

integer,parameter:: dbl=kind(1.d0)

real(dbl), pointer:: tab(:)

allocate(tab(n))

tab = 1.d0

end function func

← Cannot use “allocatable” here, must use “pointer”

Common blocks – a FORTRAN 77 construct

The use of common blocks should be avoided, and instead one should use modules.

One reason is that **common** declarations must be copied to every unit that uses them, in exactly the same order

common /params/ natoms,ndims in one unit and

common /params/ ndims,natoms in another

would cause the latter to think there are **natoms** dimensions and **ndims** atoms around,

and the compiler has no idea something is wrong.

Copying the same common declarations around is usually replaced by an **include** statement.

In this case the common blocks are in a file, say common.f, and that file is always included in a program unit if necessary as shown below:

In file common.f

common /params/ natoms, ndims

common /lattice/ nx,ny,nz

In the program that uses the common variables

Program test

include 'common.f'

...

This is almost ok, but the **include** statement is not part of the FORTRAN 77 standard, so the program may not be portable.

In practise however, this is not a serious problem, and Fortran 90 has include as standard.

All variables in a common block become accessible to all units having the same common block.

How to replace a common block with a module

FORTRAN 77 style

```
C
C  MAIN PROGRAM
C
IMPLICIT REAL*8 (a-h,o-z)
COMMON /data/ tab1(4),tab2(5)
PRINT*,'This is the main routine and it just fills two tables
$   in a subroutine and print them out'
CALL fill
PRINT*,'tab1'
DO 50 i= 1, 4
50  WRITE(*,'(5f7.3)') tab1(i)
    PRINT*,'tab2'
    DO 60 i= 1, 4
60  WRITE(*,'(5f7.3)') tab2(i)
    END
C
C  Fills tables tab1 and tab2
C
SUBROUTINE fill
IMPLICIT REAL*8 (a-h,o-z)
COMMON /data/ tab1(4),tab2(5)
DO 5 I = 1, 5
    IF(i .LE. 4) tab1(i) = 1.d0*i**2
5   tab2(i) = 1.d0*i
RETURN
END
```

FORTRAN 90/95 style

```
module data
  real(kind(1.d0)):: tab1(4),tab2(5)
end module data
!
!  MAIN PROGRAM
!
program main
  use data
  implicit none
  print*,'This is the main routine and it just fills two tables&
    & in a subroutine and print them out'
  call fill
  print*,'tab1'
  write(*,'(f7.3)') tab1
  print*,'tab2'
  write(*,'(f7.3)') tab2
end program main
!
!  Fills tables tab1 and tab2
!
subroutine fill
  use data
  implicit none
  integer:: i
  do i = 1, 5
    if(i <= 4) tab1(i) = dble(i**2)
    tab2(i) = dble(i)
  end do
end subroutine fill
```

Generic functions and subroutines: Overloading function names

Let's assume you have **N** types of data and you should do similar operations on any of them.

To begin with, you need **N** different subroutines, called **oper_type1, oper_type2 ... oper_typeN**.

In the program, you would have to call the correct subroutine for the type of data at hand.

This is not what you want to do. For example, if you want to compute the exponential of real numbers or complex numbers, you don't want to call the function “exp_real” and “exp_complex”, just “exp”. But you know already

that the intrinsic **exp()** *does* all this, so there must be a way to define a **generic subroutine oper**, and use that instead of **oper_type1, oper_type2 ... oper_typeN**. This shows how to do it.

module operations

interface **oper**

module procedure oper_type1, oper_type2, oper_type3

end interface

contains

! define operation for each data type

subroutine oper_type1(a)

declare here a to be of type 1

...

end subroutine oper_type1

subroutine oper_type2(a)

declare here a to be of type 2

...

end subroutine oper_type2

etc.

end module operations

Program main

use operations

declare here x to be of some the N types

declare here y to be of some the N types

call **oper(x)**

call **oper(y)**

end program main

You use the generic name **oper** for the operations, the compiler will **automatically** pick the correct subroutine from the list **oper_type1, ... ,oper_typeN** based on the type of the argument(s), here **x** or **y**.

All the dirty, uninteresting details are hidden in the module, the program itself is clean and easy to read

Many ways to do the same thing

Since Fortran 95 and Fortran 90 are downward compatible with FORTRAN 77, there are many syntactically different ways to do practically the same thing. The only difference is outlook and portability.

I assume here, as always, that you are using a Fortran 95 compiler, *some forms are not FORTRAN 77*.

All these define a double precision real number x

double precision :: x
double precision x
real*8:: x
real*8 x
real (kind(1.d0)) :: x **recommended**
real (kind(0.d0)) :: x **same thing**
real(2) :: x **ok, but difficult to remember**

All these define an integer i (of the same kind)

integer :: i **mostly used (see below)**
integer i
integer(kind(1)) :: i
integer(4):: i

Old syntax don't allow variable to be initialised

This is ok:

integer:: i=12345

This is **WRONG**:

integer i=12345

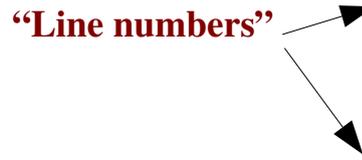
All these do-loops do the same

do 10 i = 1, 500
print*,i
10 continue

do 20 i = 1, 500
20 print*,i

do i = 1, 500 **recommended**
print*,i
end do

“Line numbers”



REMARK:

Most compilers use 32 bit integers, but some new ones may use 64 bit integers by default. The latter can break old programs. If in doubt, printing **bit_size(i)** will tell you how many bits integer **i** uses. Another useful inquiry function is **sizeof(i)**.

Loop over real points

The problem is that **real variables should not be used as loop variables**. Consider a BAD loop over real numbers x ,

```
do x=0.d0, 1.d0, 0.1d0 ! NEVER USE THIS
```

```
...
```

```
end do
```

What is the last value of x ? You might think it's $x=1$, but it's not that simple. If the addition of 0.1 to x gives a value 1.0000000000000001, to the computer this is above 1.d0 so it won't come out as $x \Rightarrow$ the biggest x *in this case* is about 0.9.

Due to this numerical glitch the loop may end at a different x in different machines, compilers or optimizations.

\Rightarrow Unpredictable behaviour – compilers warn you

Here are a few common ways to make a loop that runs over evenly spaced real number grid points $x = a, a+dx, a+2dx, \dots, b$ with n points x_i . The lowest x is always a , the last x may be b or less.

a, b and dx are given, any index

```
x = a
```

```
do
```

```
  program lines that use x
```

```
  x = x + dx
```

```
  if(x>b) exit
```

```
end do
```

a, b and dx are given, uses index i

```
n = (b-a)/dx+1
```

```
do i = 1, n
```

```
  x = a+(i-1)*dx
```

```
  program lines that use x and i
```

```
end do
```

a, b and n are given

```
dx = (b-a)/(n-1)
```

```
do i = 1, n
```

```
  x = a + (i-1)*dx
```

```
  program lines that use x
```

```
end do
```

a, dx and n are given, b unknown

```
do i = 1, n
```

```
  x = a + (i-1)*dx
```

```
  program lines that use x
```

```
end do
```

LAPACK95

LAPACK95 provides Fortran 95 style use of FORTRAN 77 routines in LAPACK/BLAS

Interfaces make possible to have : - Generic names to single and double precision routines
- Optional arguments

Drivers allocate work space and make calls to the F77 subroutines that do the actual computation

Why use LAPACK95?

- Your routines don't get messy with work space allocations/deallocations, those are now done in the driver routine you never need to see
- You don't have to worry about using a routine of wrong precision, the compiler can check for it
- The beauty of optional arguments: just leave out things you don't want.

For example, instead of calling the F77 routine with an argument that tells you *don't* want eigenvectors, you just call a generic f95 subroutine where the eigenvector argument is not present:

Example: Diagonalise a tridiagonal matrix, diagonal elements in table **d**, off-diagonal elements in table **e**

call **la_stev(d,e)** computes only the eigenvalues, returned in table **d(:)**

call **la_stev(d,e,z)** computes eigenvalues (in table **d(:)**) and eigenvectors (in table **z(:,:)**)

call **la_stev(d,e,z,info)** same as above, but with integer **info** returned

(if **info** is not present and the program fails, it will terminate and print an error message)

- The LAPACK95 drivers have a better chance to be the same in the future, even if someone changes the computing routines in LAPACK

The generic LAPACK95 routines are called **la_XXXX**, the “**la_**” is there to distinguish from LAPACK names

LAPACK naming scheme

Taken from www.netlib.org

The name of each LAPACK95 routine has been made as similar as possible to its name in LAPACK. All driver and computational routines have names of the form LA_YYZZZ, where for some routines the 8 character is blank. The two letters YY indicate the type of matrix (or of the most significant matrix). Most of these two-letter codes apply to both real and complex matrices; a few apply specifically to one or the other, as indicated in Table 2.1.

Table 2.1: Matrix types in the LAPACK naming scheme

GB	general band
GE	general (i.e., unsymmetric, in some cases rectangular)
GG	general matrices, generalized problem (i.e., a pair of general matrices)
GT	general tridiagonal
HB	(complex) Hermitian band
HE	(complex) Hermitian
HP	(complex) Hermitian, packed storage
PB	symmetric or Hermitian positive definite band
PO	symmetric or Hermitian positive definite
PP	symmetric or Hermitian positive definite, packed storage
PT	symmetric or Hermitian positive definite tridiagonal
SB	(real) symmetric band
SP	symmetric, packed storage
ST	(real) symmetric tridiagonal
SY	symmetric

When we wish to refer to a class of routines that perform the same function on different types of matrices, we replace the two letters by ``yy''. Thus LA_yySV refers to all the simple driver routines for systems of linear equations that are listed in Table 2.2. The last three letters ZZZ indicate the computation performed.

LAPACK95 cont'd : Example of an interface

Advanced feature

```
INTERFACE LA_GEGS
  SUBROUTINE DGEGBS_F95( A, B, ALPHAR, ALPHAI, BETA, VSL, VSR, &
    &          INFO )
    USE LA_PRECISION, ONLY: WP => DP
    INTEGER, INTENT(OUT), OPTIONAL :: INFO
    REAL(WP), INTENT(INOUT) :: A(:,,:), B(:,:)
    REAL(WP), INTENT(OUT), OPTIONAL :: ALPHAR(:), ALPHAI(:), &
    &          BETA(:)
    REAL(WP), INTENT(OUT), OPTIONAL, TARGET :: VSL(:,,:), VSR(:,:)
  END SUBROUTINE DGEGBS_F95

  SUBROUTINE ZGEGS_F95( A, B, ALPHA, BETA, VSL, VSR, INFO )
    USE LA_PRECISION, ONLY: WP => DP
    INTEGER, INTENT(OUT), OPTIONAL :: INFO
    COMPLEX(WP), INTENT(INOUT) :: A(:,,:), B(:,:)
    COMPLEX(WP), INTENT(OUT), OPTIONAL :: ALPHA(:), BETA(:)
    COMPLEX(WP), INTENT(OUT), OPTIONAL, TARGET :: VSL(:,,:), &
    &          VSR(:,:)
  END SUBROUTINE ZGEGS_F95
END INTERFACE
```

Generic name for two routines:

- **Double precision routine**
- **Double Complex routine**

DP stands for "Double Precision"

Optional arguments: may be present or not

FORALL: a new f95 program structure

FORALL is a loop construct for doing things that can be done in parallel, meaning that the order of doing things is irrelevant. Even in a single processor machines a good compiler can take advantage of the many registers in the processors to do many tasks at once. This was the idea. In reality compilers are much better at optimising DO's than FORALL's, so you probably find out that a FORALL is either as fast or slower than DO - sigh.

program test

implicit none

integer, parameter:: n=7000

integer:: i

real(kind(1.d0)):: A(n)

forall (i=1:n)

A(i) = i2**

end forall

end program test

We can use forall, because the values of elements A(i) are independent.

For example, it doesn't matter whether we set $A(1000) = 10002$ before or after we set $A(2123) = 2123**2$.**

The forall loop above can also be written on a single line:

forall(i=1:n) A(i)

=i2**

SELECT CASE structure

A multiple-choice branching of the program flow. Often more readable than a sequence of IF tests.

```
integer:: i
```

```
...
```

```
select case(i)
```

```
  case (:-1)
```

```
    print*, 'i is negative'
```

```
  case (0)
```

```
    print*, 'i is zero'
```

```
  case(1:)
```

```
    print*, 'i is positive'
```

```
  case default
```

```
    print*, 'i is probably NaN'
```

```
end select
```

```
...
```

case (:-1) means all numbers ..., -4,-3,-2,-1



If none of the cases before matched



Numerical Recipes *by Press, Teukolsky, Vetterling and Flannery, see www.nr.com*

Fortran 77, C, C++ and Fortran 90

- **A collection of subroutines for numerical tasks : eigenvalues, FFT, random numbers *etc.***
- **Very well known among the numerical community**
- **The efficiency of the programs divides opinions:**
They get the job done, but are not guaranteed to give the best method for the specific task
- **The manual is a nice bedside reading for those interested in numerical methods: some of the NR books are also free online**
- **To obtain a copy of the subroutines buy them with the book – many machines have them installed already**

FFTW *see www.fftw.org*

Written in C, Fortran wrappers

- **“Fastest Fourier Transform in the West”**
- **Some vendor-tuned FFT's may be faster, but FFTW is at least almost as fast and portable**
- **If your program spends appreciable time in doing FFT's, use FFTW**
If not, use any simple FFT, like the one in Numerical Recipes
- **Older version 2.x has different API than the newer 3.x**

Preprocessing:

The difference between program.f90 and program.F90

- **program.f** : fixed form, things have to be in specific columns
a FORTRAN 77 routine
- **program.F** : The compiler is instructed to do preprocessing to make a .f file out of the .F file (visible or not)

- **program.f90** : free form (no column specific rules, but keep them readable and use intendation!)
a Fortran 90 routine
- **program.F90** : The compiler is instructed to do preprocessing to make a .f90 file out of the .F90 file

Preprocessing is a way to make choices at compile time

For example, you may want to use the FFT provided by the ESSL library in IBM and the FFTW library in other machines

Example: file program.F90

```
program test
#ifdef PARA
    print*,'here I will put my parallel code once I learn how ...'
#else
    print*,'this is my serial machine code'
#endif
    print*,'this part is done in both serial and parallel cases'
end program test
```

Preprocessor directive

Compile the example using

a) `f90 program.F90 -o goserial`

and

b) `f90 -DPARA program.F90 -o gopara`

and you have two different binary programs, goserial and gopara

Try also `cpp program.F90` and `cpp -DPARA program.F90`

Don't put preprocessor directives in a file with suffix .f90:

fortcom: Warning: badexample.f90, line 3: Bad # preprocessor line

```
#ifdef PARA
```

```
_^
```


Numerical constants, cont'd

If you have a hard-wired constant always define it as a **parameter**:

- You don't risk changing a parameter by mistake
- If you decide to change, say, 110 to 220 you don't want to edit all files that had that 110. There is a risk that you miss one or more occurrences of that “110”.

Never ever do like this:

```
program test
  implicit none
  integer:: i,index(100)
  do i = 1, 100
    index(i)= i
  end do
  call rotateindex(index)
  write(12,*) index
end program test
subroutine rotateindex(index)
  implicit none
  integer:: index(100),i
  i = index(1)
  index(1:99) = index(2:100)
  index(100) = i
end subroutine rotateindex
```

Change only this
and you are done

Do it like this:

```
program test
  implicit none
  integer, parameter:: n = 100
  integer:: i,index(n)
  do i = 1, n
    index(i)= i
  end do
  call rotateindex(index,n)
  write(12,*) index
end program test
subroutine rotateindex(index,n)
  implicit none
  integer:: n, index(n),i
  i = index(1)
  index(1:n-1) = index(2:n)
  index(n) = i
end subroutine rotateindex
```

If you wanted to change 100 to 101,
then even in this short program
you would have to edit all these 6 points
to get it working again.

Remark 1: You cannot replace all “100” by
“101” automatically, it's not enough
– and often it's too much !

Remark 2: The subroutine works only for a
100 element input: a very dull idea.

This is just an example. If you want to rotate indices use the built-in function “shift”.

The scope of variables

When is a variable visible to a function or a subroutine?

This won't work:

Program test

```
implicit none
integer:: i,k=100
do i = 1, 5
  call addone
  print*,k
end do
end program test
```

subroutine addone

```
integer:: k
k = k + 1
end subroutine addone
```

**Variable `k` in the main program
and the variable `k` in the subroutine addone
are not the same. In fact, the `k` in the subroutine is
completely arbitrary!**

**Add here “print*,k”
to see what the subroutine thinks `k` is**

Output:

```
100
100
100
100
100
```

There is no way this subroutine, which is outside the main program, receives any information about the variable `k`.

The scope of variables, cont'd

When is a variable visible to a function or a subroutine (or to the calling program unit)?

This works:

```
program test
  implicit none
  integer:: i,k=100
  do i = 1, 5
    call addone
    print*,k
  end do
contains
  subroutine addone
    k = k + 1
  end subroutine addone
end program test
```

← This k is visible to the subroutine, because ...

← ... the subroutine is contained in the main program and sees all its variables.

Output:

101
102
103
104
105

And how easy it is to break it again!

```
program test
  implicit none
  integer:: i,k=100
  do i = 1, 5
    call addone
    print*,k
  end do
contains
  subroutine addone
    integer:: k
    k = k + 1
  end subroutine addone
end program test
```

BAD! →

Why bad? The added line
integer:: k
defined another variable k;
this new k is an internal
variable of the subroutine.

Output:

100
100
100
100
100

The scope of variables, cont'd

Private and public variables: yet another attribute a variable can have

Advanced feature

A variable may be defined to be private or public to control it's visibility to other program units.

module params

```
! real(kind(1.d0)),private:: salary,taxes  
real(kind(1.d0)):: salary,taxes
```

contains

```
function taxprocent()
```

```
salary = 13000.d0
```

```
taxes = 4210.d0
```

```
taxprocent = taxes/salary*100.d0
```

```
end function taxprocent
```

```
end module params
```

program test

```
use params
```

```
implicit none
```

```
write(*,('Tax Procent is ',f0.2)) taxprocent
```

```
()
```

```
write(*,(' Salary is ',f0.2)) salary
```

```
end program test
```

**The commented line has an additional attribute “private”;
It would state that salary and taxes are hidden from
any users of the module params.**

**In that case the main program couldn't print
out the salary (compiler gives an error message)**

**All variables in a module are by default public
(visible to anyone who uses the module)**

**You can change the default from
public to private:**

```
Module something
```

```
private
```

```
real(kind(1.d0)):: ...
```

```
integer::...
```

```
end module something
```

Or you can specify the visibility of each variable:

```
Module something
```

```
real(kind(1.d0)), public:: ...
```

```
integer, private:: ....
```

```
end module something
```

Good programming habits

- A good main program is a collection of subroutine calls.

The main program is a leader, it sees over that the principal tasks get done in correct order.

=> Don't put any lengthy code segments to the main program.

```
program main
  use parameters
  implicit none
  integer:: i
  call initialize
  do i = 1, Niter
    call iterate
    if(i>Ntherm) call measure
  end do
  call errorestimation
  call saveresults
end program
```

- Always INDENT programs. Not too much, not too little. Nobody likes to read unindented program.
- Choose descriptive, *short* names. Name loops that end far away from start
- Keep formulas close to their **mathematical form**. That form has been around for decades and shows *what* you compute.

NEVER WRITE

```
do i = 1, n
  y(i) = sin(i+j/20*n+4)*cos(i+j/20*n+4)*exp( i+j/20*n+4)
end do
```

BUT WRITE

```
do i = 1, n
  x = i+j/20*n+4
  y(i) = sin(x)*cos(x)*exp(x)
end do
```

- Comment things that are not obvious.
- Don't comment everything. Your program is not a Fortran manual. If the reader doesn't know Fortran, it's not *your* fault.

Using only some variables from a module

Syntax:

```
use module_name, only: list_of_variables
```

Example:

```
module params
  integer, parameter:: dp=kind(1.d0)
  real(dp),parameter:: pi=3.141592653589793238d0
  real(dp):: x,y,z
end module params
```

```
program test
  use params, only: x
  implicit none
  x = 10.d0
  print*,x
end program test
```

← The restriction **only** means we take from the module params only **x**. If we wanted, we could define a variable **y** within the main program without touching the other **y** in the module.

The constants **dp** and **pi** and variables **y** and **z** are not made known to the main program.

Terminology

About arrays

Rank = The number of array dimensions

`integer,allocatable:: index(:,::,)` array index has rank 3

Shape = Tells what indices does the array cover

`integer :: index(5,6,2)` array index has shape (5,6,2)

Explicit shape array = The upper bound of each dimension is specified

`integer:: index(5,6,2)` or `integer:: index(k,l,m)`

Simple, but not very versatile; sometimes require a lot of checking by you

Assumed size array = A Fortran 77 relict. The upper bound of some dimension is not specified

`integer:: index(5,6,*)`

Avoid, use assumed shape arrays instead

Assumed shape array = The rank is specified, but not the shape

`integer:: index(:,::,)` can have any shape, such as `index(5,6,2)` or `index(1,2,2)` or ...

Use frequently! Put these in modules or write an explicit interface.

Remember also, that all lower bounds are 1 by default, unless told otherwise.

About functions and subroutines

Dummy argument = Argument in the *declaration* of a function or a subroutine, a temporary name.

`real(kind(1.d0)) function cube(x)`

here `x` is a dummy argument

`subroutine mysub(input,output)`

here `input` and `output` are dummy arguments

Actual argument = Argument in the *call* of a function or a subroutine, the name of the thing that is actually sent to the function or the subroutine to operate on

`x3 = cube(x)`

here `x` is the actual argument

`call subroutine mysub(testinput, outcome)`

here `testinput` and `outcome` are actual arguments

Command line arguments

Sometimes writing a namelist or a parameter file is too clumsy. For example, if you have written a program “sort” that sorts a file containing numerical data to ascending order according to some column, then a natural way to use it is from a command line

```
> sort file column
```

Below is a short example how to read in the command line – without any testing of error conditions to keep it basic.

Program main

```
implicit none
```

```
character(100):: buffer
```

```
character(25):: filename
```

```
integer:: column
```

← A long string to hold the command line

← Two command line arguments in this example

```
call getarg(1,buffer)
```

```
read(buffer,*) filename
```

```
call getarg(2,buffer)
```

```
read(buffer,*) column
```

```
...
```

```
end program main
```

← Notice how one reads from a character string “buffer” !

Internal read and write

One can read and write to and from other “devices” than the keyboard and files. One can also read from and write to **internal** sources. The example shows how to write to a character string.

```
program main
  implicit none
  integer:: i
  character(50):: s
  print*, 'give an integer (not too long)'
  read(*,*) i
  write(s, '(i0)') i
  s = 'prefix.//trim(s)//'.suffix'
  print*, 'Making the character string "prefix.i-value.suffix"'
  print*, 'The result is ', trim(s), ' as it should'
end program main
```

s is 50 characters long => lots of blanks to remove

```
>a.out
  give an integer (not too long)
325636
  Making the character string "prefix.i-
value.suffix"
  The result is prefix.325636.suffix as it should.
```

Strict “i0” format together with trim(s) removed blanks that would otherwise spoil the output

string1//string2 use “//” to glue strings together

trim(string) removes trailing blanks from string

Few words about parallel programming

- **HPF = High Performance Fortran**

A collection of parallelisation directives on top of Fortran 90/95

Shows in the program as lines starting

!HPF\$

Online course: <http://www.liv.ac.uk/HPC/HPFpage.html>

<http://dacnet.rice.edu/Depts/CRPC/HPFF/index.cfm>

- **MPI, especially Open**

MPI = Message Passing Interface

MPI

Open MPI is a free implementation of MPI

Open MPI Represents the merger between three well-known MPI implementations:

FT-MPI from the University of Tennessee

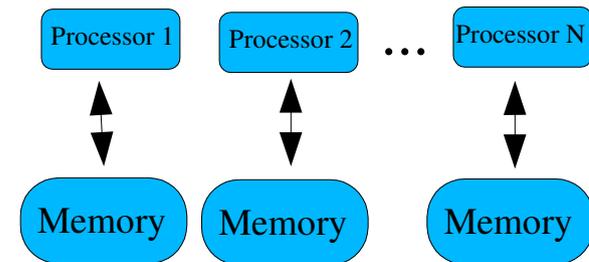
LA-MPI from Los Alamos National Laboratory

LAM/MPI from Indiana University

Can be installed even on single processor machines :^)

www.open-mpi.org
www.tu-darmstadt.de/hrz/hhhr/doku/sw/pe/am106mst02.html

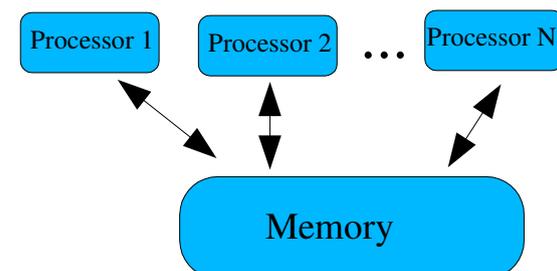
Distributed-memory systems



More about MPI on the next slide!

www.openmp.org

Shared-memory systems



- **OpenMP**

MP = Multi Processing

A shared-memory parallel programming environment

Parallelisation directives

Library routines

Shows in the program as lines starting

!\$OMP (in C they start **#pragma omp**)

MPI: *de facto* standard of parallel programming

Only choice for distributed-memory systems: All processors access their own memory; *e.g.* PC clusters

We have *nodes* processors (nodes): master is 0, slaves are *1,2,...,nodes-1*. Also called root and workers.

The parallel program calls MPI subroutines that initialise, distribute input data and collect results. The set of nodes make a group called *comm* and the node (root or a slave) where the code is running is number *rank*.

Some common MPI subroutines: (*ierr*=integer for an error, see *man mpi_XXXX* for more information on the other arguments)

mpi_init(*ierr*) Start MPI

mpi_comm_size(*comm,nodes,ierr*) How many nodes belong to the group *comm* (usually: # of processors you use)

mpi_comm_rank(*comm,rank,ierr*) Get the rank (0, 1, ... nodes-1) of the current processor

mpi_bcast(*buffer, count, datatype, master,comm, ierr*) Broadcasts a message (data) from the master to all other processes of the group

mpi_barrier(*comm,ierr*) All processors wait at the barrier for the rest to arrive

mpi_reduce(*sendbuf,recvbuf,count,datatype,op,master,comm,ierr*) Reduces values on all processes to a single value, *op* = operation (*mpi_sum*)

mpi_finalize(*ierr*) End parallel execution

----- A parallel program can be made with just the 7 subroutines given above -----

mpi_send(*buf,count,datatype,dest,tag,comm,ierr*) Basic send operation

mpi_recv(*buf, count,datatype,source,tag,comm,status,ierr*) Basic receive operation

mpi_scatter(*sendbuf,sendcount,sendtype,recvbuf,recvcount,recvtype,master,comm,ierr*) Scatter data from one task to all other tasks in a group

or mpi_scatterv(*sendbuf,sendcounts(*),displs(*),sendtype,recvbuf,recvcount,recvtype,master,comm,ierr*) Scatter data in parts

This is an excerpt from the Open MPI FAQ at www.openmpi.org, showing what problems generic names (subroutine overloading) may cause.

14. Why does compiling the Fortran 90 bindings take soooo long?

This is actually a design problem with the MPI F90 bindings themselves.

The issue is that since F90 is a strongly typed language, we have to overload each function that takes a choice buffer with a typed buffer.

For example, MPI_SEND has many different overloaded versions -- one for each type of the user buffer.

Specifically, there is an MPI_SEND that has the following types for the first argument:

logical*1, logical*2, logical*4, logical*8, logical*16 (if supported)
integer*1, integer*2, integer*4, integer*8, integer*16 (if supported)
real*4, real*8, real*16 (if supported)
complex*8, complex*16, complex*32 (if supported)
character

On the surface, this is 17 bindings for MPI_SEND. Multiply this by every MPI function that takes a choice buffer (50) and you 850 overloaded functions. However, the problem gets worse -- for each type, we also have to overload for each array dimension that needs to be supported. Fortran allows up to 7 dimensional arrays, so this becomes $(17 \times 7) = 119$ versions of every MPI function that has a choice buffer argument. This makes $(17 \times 7 \times 50) = 5,950$ MPI interface functions.

To make matters even worse, consider the ~25 MPI functions that take 2 choice buffers. Functions have to be provided for all possible combinations of types. This then becomes exponential -- the total number of interface functions balloons up to 6.8M.

Additionally, F90 modules must all have their functions in a single source file. Hence, all 6.8M functions must be in one .f90 file and compiled as a single unit (currently, no F90 compiler that we are aware of can handle 6.8M interface functions in a single module).

To limit this problem, Open MPI, by default, does not generate interface functions for any of the 2-buffer MPI functions. Additionally, we limit the maximum number of supported dimensions to 4 (instead of 7). This means that we're generating $(17 \times 4 \times 50) = 3,400$ interface functions in a single F90 module. So it's far smaller than 6.8M functions, but it's still quite a lot.

This is what makes compiling the F90 module take so long.

OpenMP

Shared-memory systems: all processors access the same memory; e.g. Origin 2000 , IBM PS

- Ease of use
- Incremental parallelization
- Easy speedup on desktops (dual core machines, threads) - you can start low
- Scales well on multiprocessor machines - you can get high

Intel f90 compiler (and many other) supports OpenMP directives. Since the directives look like ! \$omp xxxxxxxx a compilation using

```
ifort program.f90
```

 (this fails if there are calls to openMP library routines)

produces ordinary (serial) code, while compilation using

```
ifort -openmp program.f90
```

produces parallel code following the directives. => You have *both* a serial and a parallel version of your program.

[Excerpt from the openMp sample routine ; molecular dynamics www.openmp.org/drupal/samples/md.html](http://www.openmp.org/drupal/samples/md.html)

By Bill Magro, Kuck and Associates, Inc. (KAI), 1998

This loop updates the positions and velocities of the particles using the velocity Verlet algorithm

```
...
! The time integration is fully parallel
!$omp parallel do
!$omp& default(shared)
!$omp& private(i,j)
do i = 1,np
do j = 1,nd
pos(j,i) = pos(j,i) + vel(j,i)*dt + 0.5*dt*dt*a(j,i)
vel(j,i) = vel(j,i) + 0.5*dt*(f(j,i)*rmass + a(j,i))
a(j,i) = f(j,i)*rmass
enddo
enddo
!$omp end parallel do
...
```

OpenMP directives for parallelization of the loop

```
> setenv OMP_NUM_THREADS 8
```

 (run with 8 threads/processors)

```
> ifort -openmp md.f
```

```
md.f(92) : (col. 7) remark: OpenMP DEFINED LOOP WAS PARALLELIZED.
```

```
md.f(209) : (col. 7) remark: OpenMP DEFINED LOOP WAS PARALLELIZED.
```

(the other parallelized loop was force calculation)

OpenMP timings

program md.f (openMP sample program)

lilli : SGI Altix 3000 Intel Itanium 2, 900MHz/1.5MB L3 Cache

origin: SGI Origin 3800 R12000 400MHZ

compiled: lilli ifort -fast -openmp md.f

origin f90 -Ofast=ip27 -mips4 -64 -mp md.f

(f90 = MipsPro)

Execution times in seconds

#proc	lilli	origin	Excellent scaling !
1	57	72	
2	35	36	
4	18	19	
8	10	9	
16	6	5	

Complex arithmetics: application to Mandelbrot set

program mandel

implicit none

integer, parameter:: dp=kind(1.d0),ix=10000 ! Max # of iterations

! resolution d, box corners [x0,y0] and [x1,y1]

real(dp),parameter::d=1.d-2,x0=-1.5d0,x1=0.5d0,y0=-1.d0,y1=1.d0

integer:: i

complex(dp):: cx,c

real(dp):: x,y

x = x0

do

 y = y0

 do

 c = cmplx(x,y,dp)

 cx = c

 do i = 1, ix

 cx = cx**2+c

 if(abs(cx)>2) exit

 end do

 write(12,'(2f8.3,i10)') c,i

 y = y + d

 if(y>y1) exit

 end do

 write(12,*)

 x = x + d

 if(x>x1) exit

end do

end program mandel

Plotting using gnuplot:

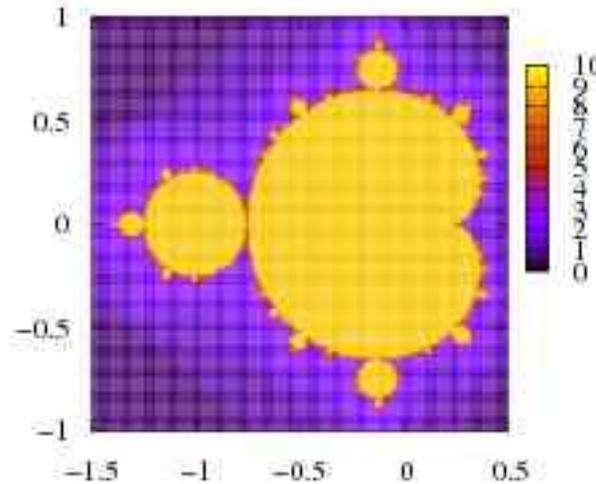
>gnuplot

gnuplot > set size ratio 1

gnuplot > set pm3d map

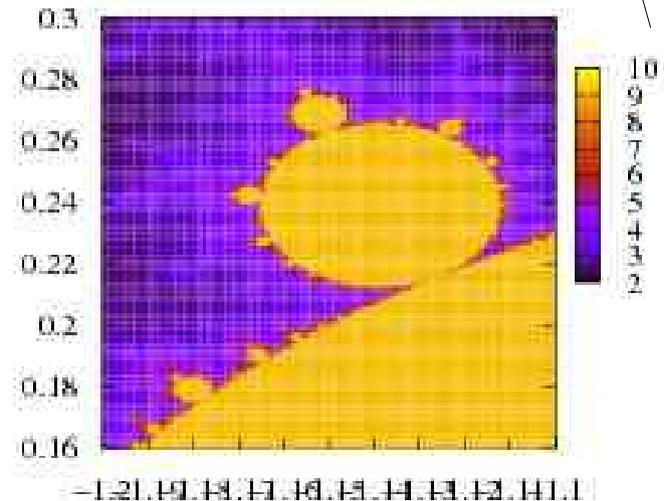
gnuplot > unset key

gnuplot > sp 'fort.12' u 1:2:(log(\$3))



$d=1.d-2,x0=-1.5d0,x1=0.5d0,y0=-1.d0,y1=1.d0$

$d=5.d-4,x0=-1.20d0,x1=-1.1d0,y0=0.16d0,y1=0.3d0$



Notice how close the program is to the mathematical algorithm!

Algorithm

$cx=c$

$cx=cx^2+c$ iterate max ix iterations

if $|cx|>2$ the point c is definitely outside the set

store how many iteration it takes to get out, i

Mandelbrot set : A parallel version

Intel f90 compiler (version 9) can autoparallelize code via openMP. The previous Mandelbrot program won't autoparallelize (fastest way to see it is just to try), so let's tweak it a bit.

There are two problems to fix:

- IO from innermost loop : which processor/thread is supposed to do that?
- The outer do-loops had no loop index, instead there were exit tests **if(x>x1) exit**

If the x-loop is parallelized each process must exit the loop at a different x, which is not x1.

```
program mandel
  implicit none
  integer, parameter:: dp=kind(1.d0),imax=10000
  real(dp),parameter:: d=1.d-2,x0=-1.5d0,x1=0.5d0,y0=-1.d0,y1=1.d0
  integer:: i,ix,iy,nx,ny
  complex(dp):: cx,c
  real(dp):: x,y ,t0,t1
  integer,allocatable:: set(:,:)
  nx = (x1-x0)/d + 1
  ny = (y1-y0)/d + 1
  allocate(set(nx,ny))
  do ix = 1, nx
    x = x0+(ix-1)*d
    do iy = 1, ny
      y = y0+(iy-1)*d
      c = cmplx(x,y,dp)
      cx = c
      do i = 1, imax
        cx = cx**2+c
        if(abs(cx)>2) exit
      end do
      set(ix,iy)= i
    end do
  end do
  ... cut away: output result table "set" to file...
end program mandel
```

Clearly independent loop variables ix and iy

No file output
Table set can be filled in arbitrary order

Output of intel ifort version 9.0

Optimize
Autoparallelize
Print out a report about autoparallelization

```
> ifort -O3 -parallel -par_report3 mandel.para.f90
procedure: mandel
serial loop: line 26: not a parallel candidate due to the loop being lexically discontinuous ← The i loop
serial loop: line 22: not a parallel candidate due to insufficient work ← The iy loop
serial loop: line 38: not a parallel candidate due to statement at line 40 ← The write statement
serial loop: line 36: not a parallel candidate due to statement at line 40
mandel.para.f90(20) : (col. 3) remark: LOOP WAS AUTO-PARALLELIZED.
parallel loop: line 20 ← The ix loop
  shared   : { "I.SI32.var$16_dv_template.dim_info.lower_bound.betype.0.0"
"I.SI32.var$16_dv_template.dim_info.spacing.betype.0.0" "P.P32.var$16_dv_template.addr_a0.betype.0.0" }
  private  : { "ix" "x" "iy" "y" "c" "cx" "i" }
  first priv.: { }
  reductions : { }
```

**Result is a working parallel binary, just set the environment variable
OMP_NUM_THREADS to the number of processors/threads and run.**