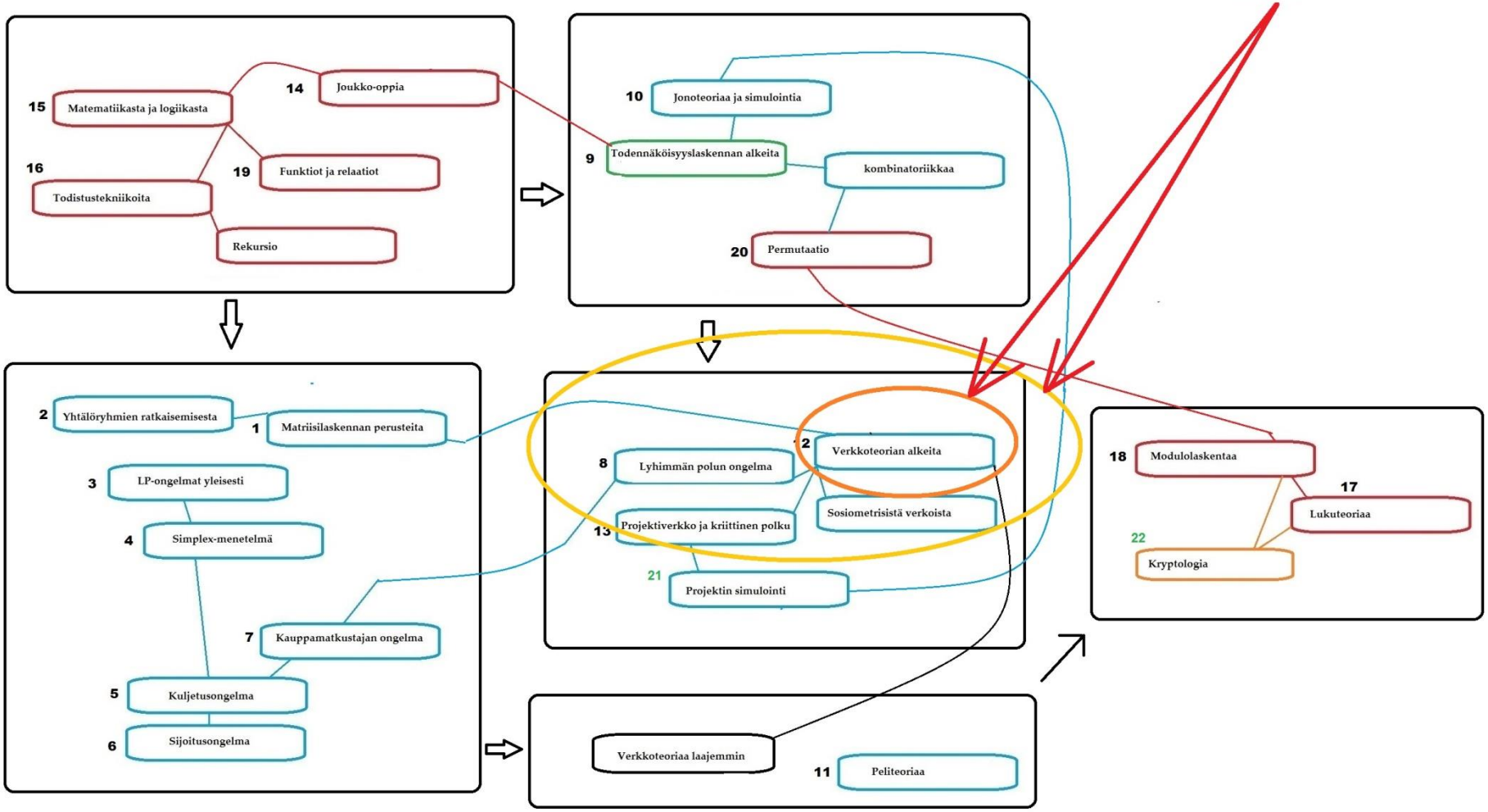
ERÄS MAHDOLLINEN JAOTTELU:

TIEP TUHAT: 2017



VERKOISTA TIIVISTETYSTI

VERKKOTEORIA = GRAAFITEORIA

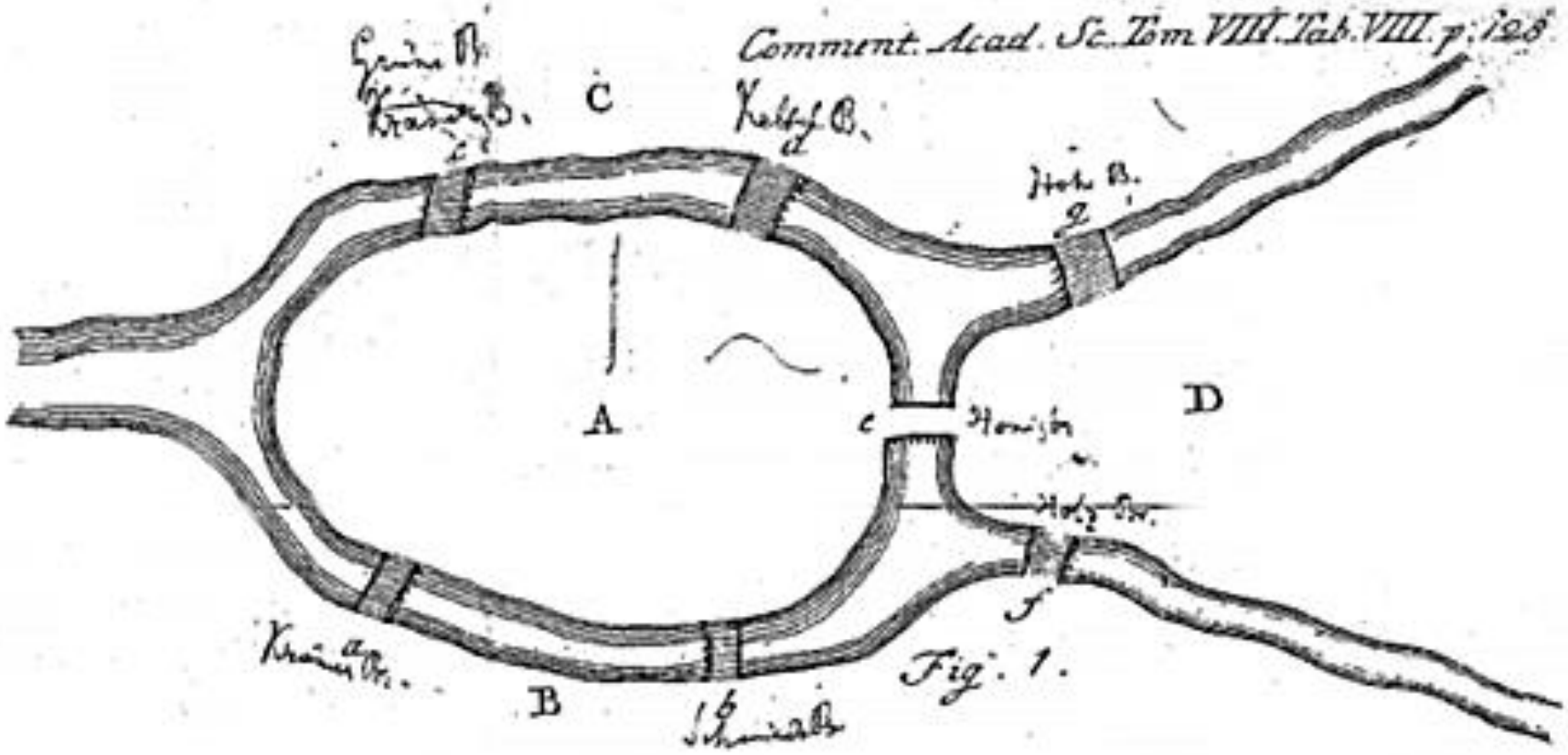
SOLMU = KÄRKI = PISTE

VÄLI = KAARI = NUOLI = LINKKI

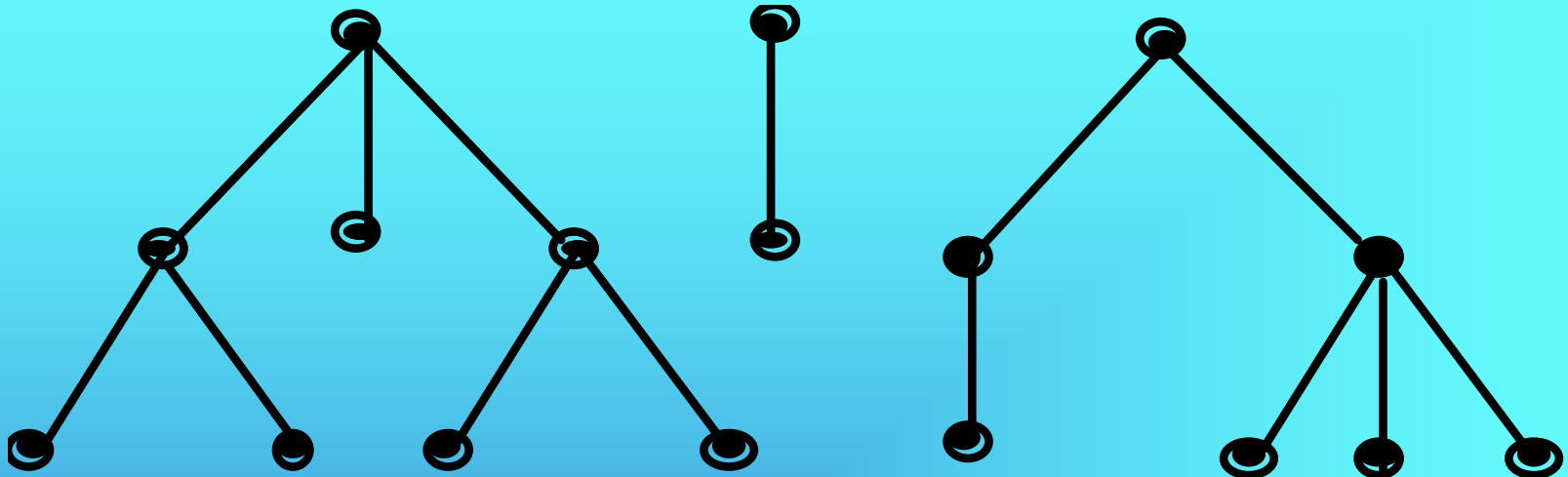
“RATSUN POLKU” = RATSUN KÄVELY”

J.N.E…

Netspace


Comment. Acad. Sc. Tom VIII. Tab. VIII. p. 128.
C
D
A
B
Fig. 1.

TASOVERKKO
ALIVERKKO
SOLMUN ASTE
LEIKKAUSSOLMU
SOSIOGRAMMI
EULERIN KAAVA
4CC
LYHIN POLKU
KRIITTINEN POLKU



PUU, METSÄ JA LEHDET

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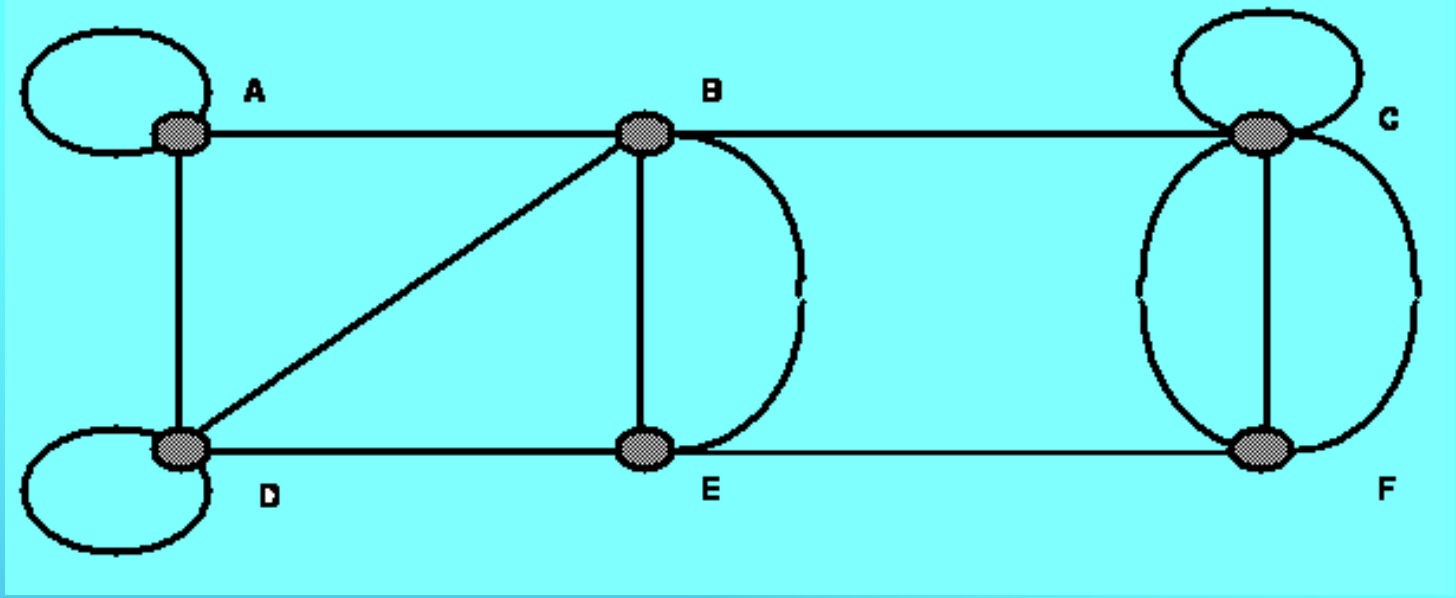
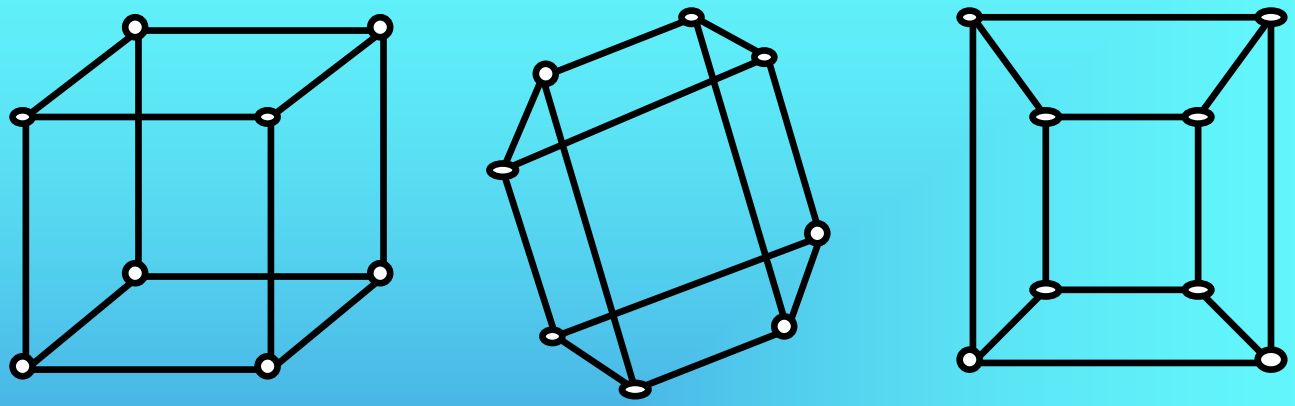
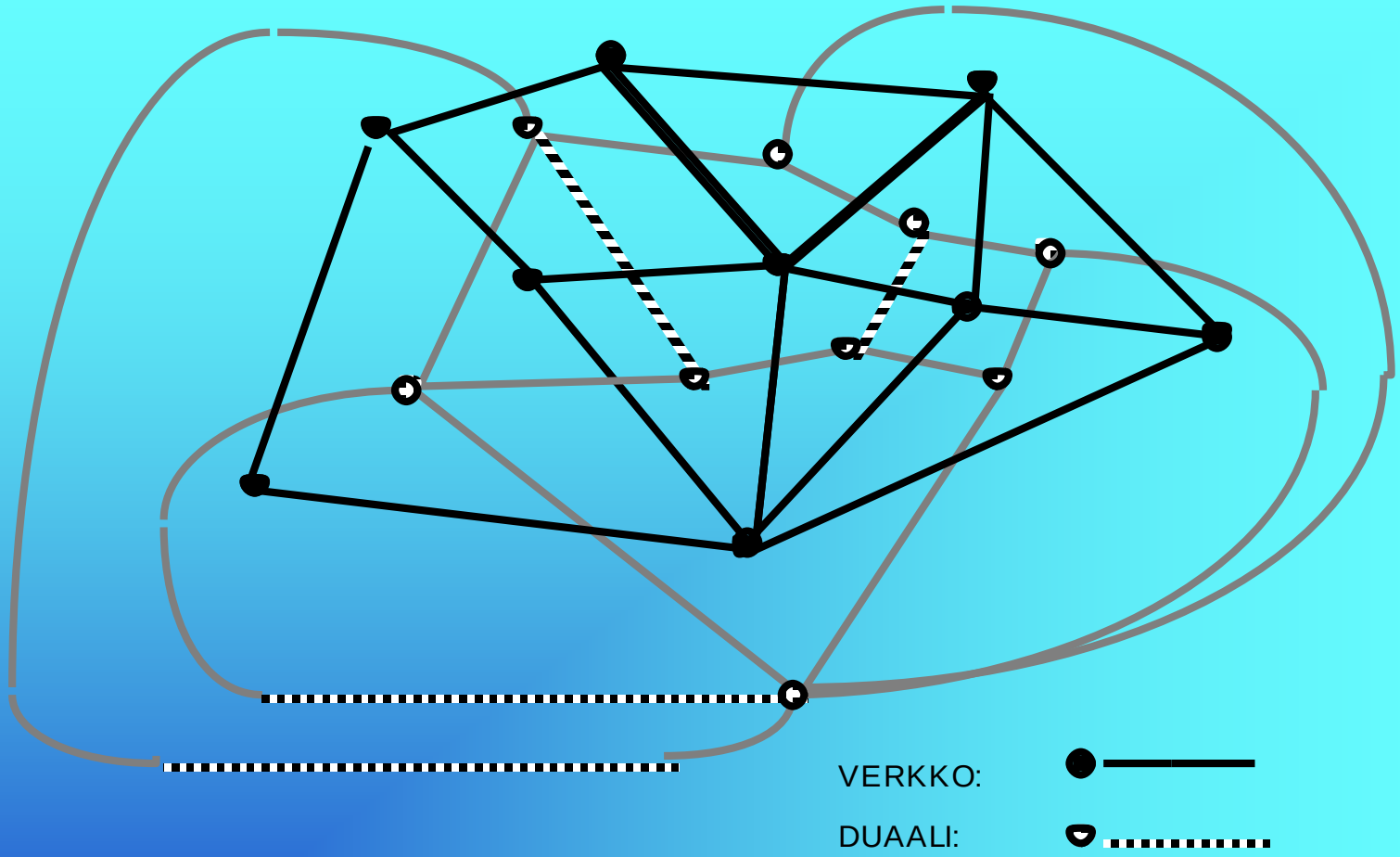


SOLMU, VÄLI, ALUE

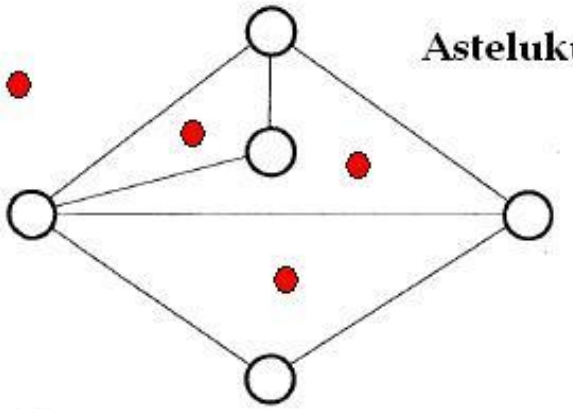
Pseudo graph



ISOMORFISET VERKOT

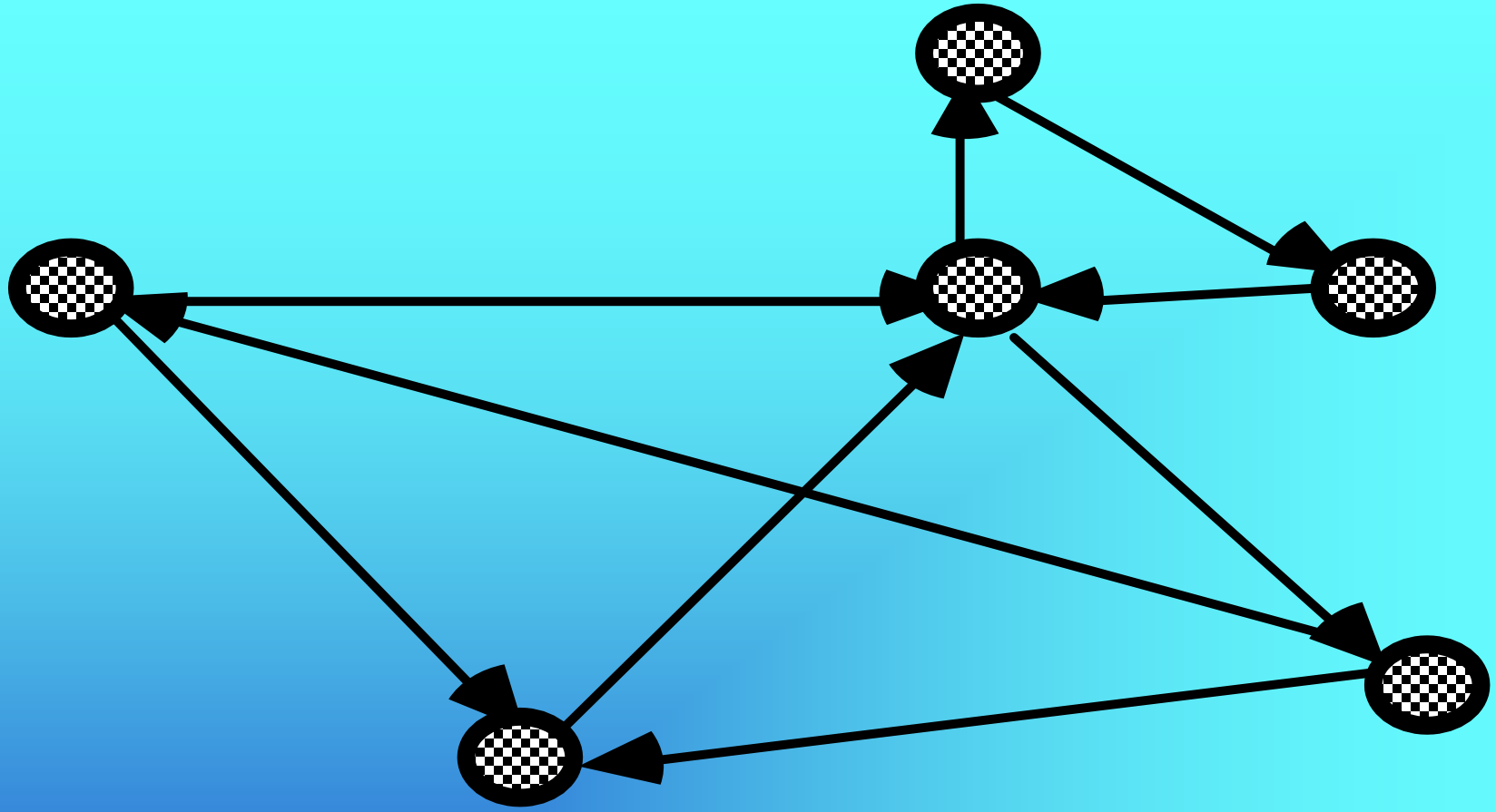


VERKON DUAALI



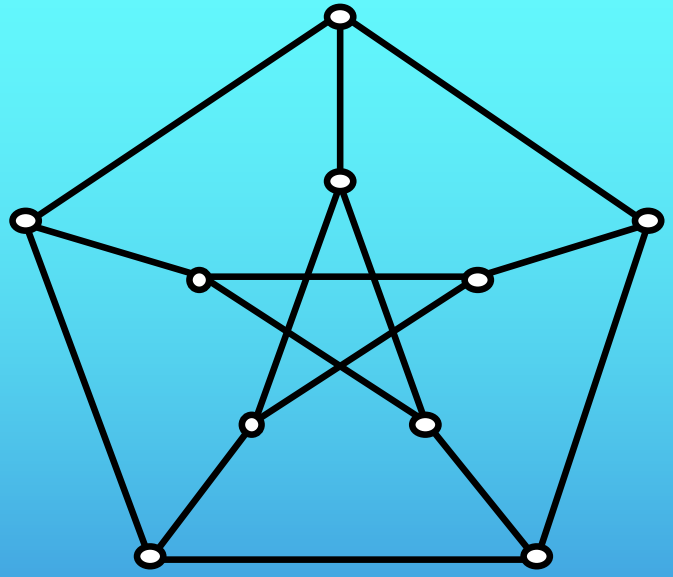
Astelukua

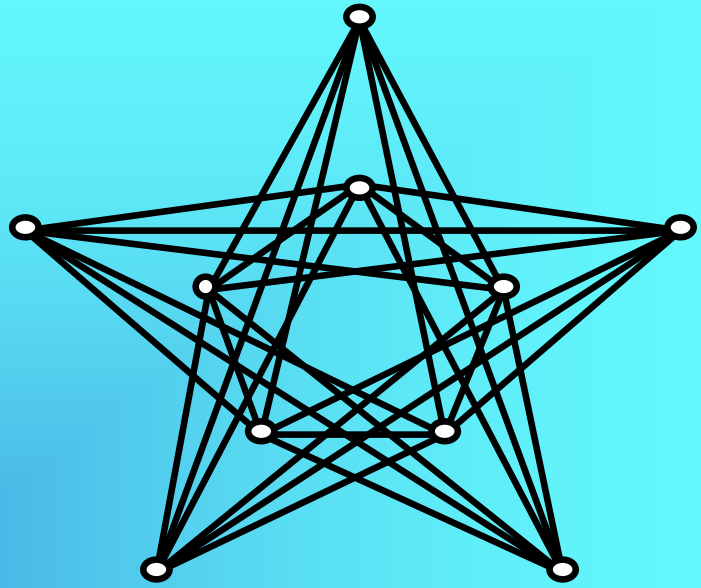
5	4	3
	2	2
1		3
1		3



SUUNNATTU VERKKO

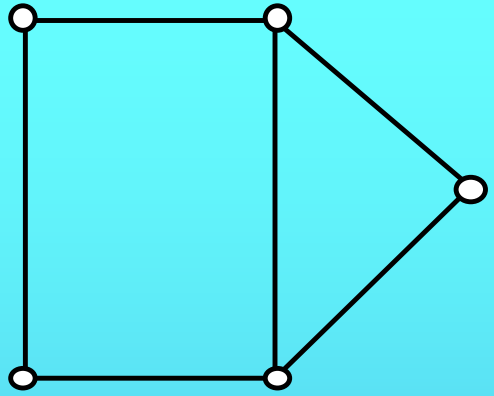
Digraph

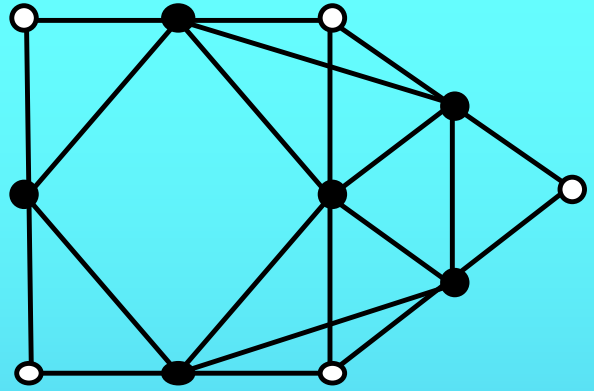


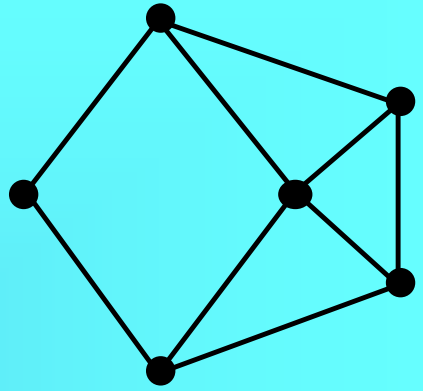


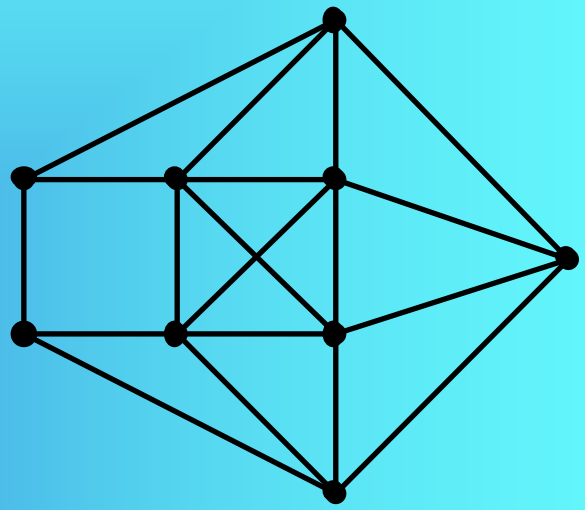
VERKON KOMPLEMENTTI

Petersen



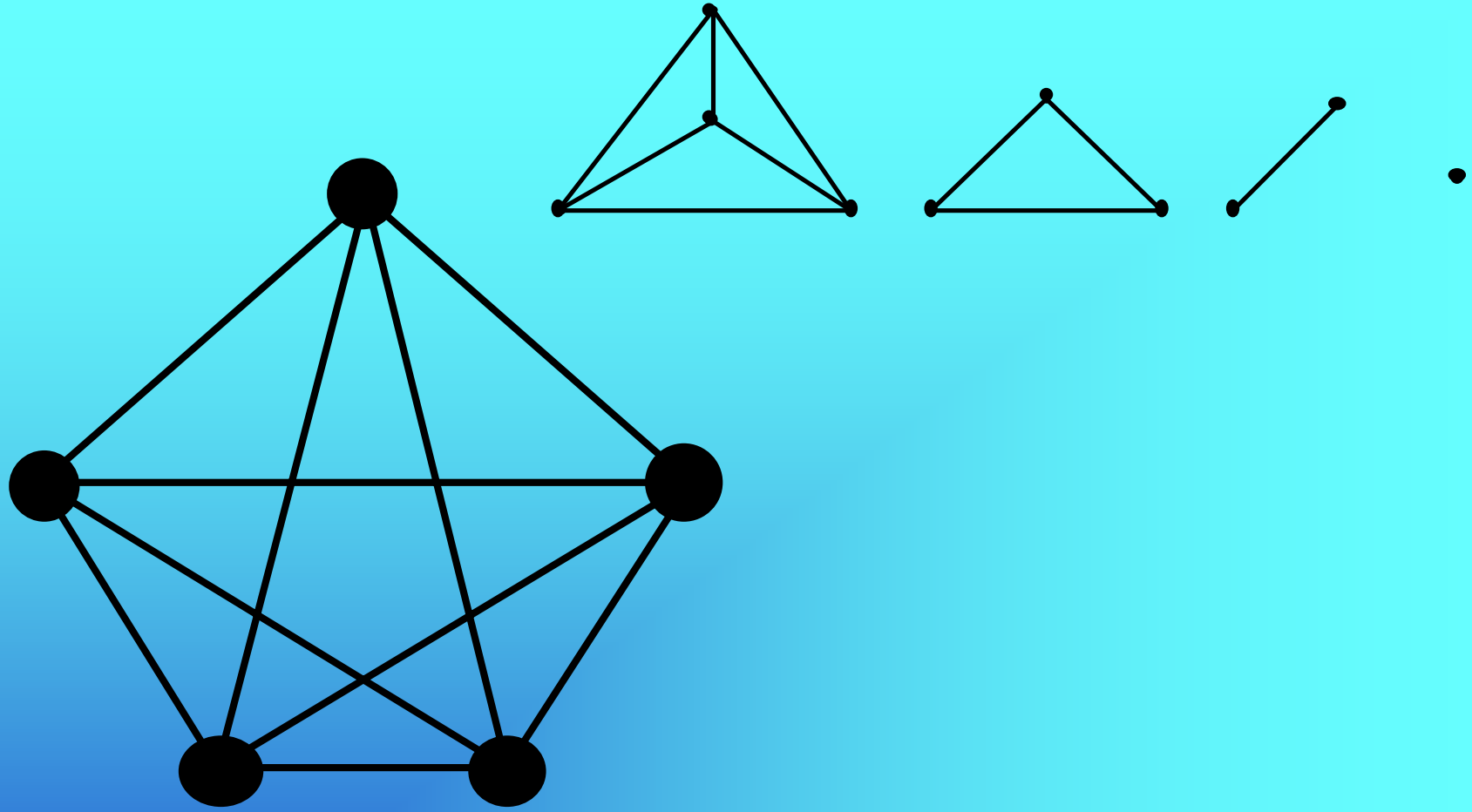






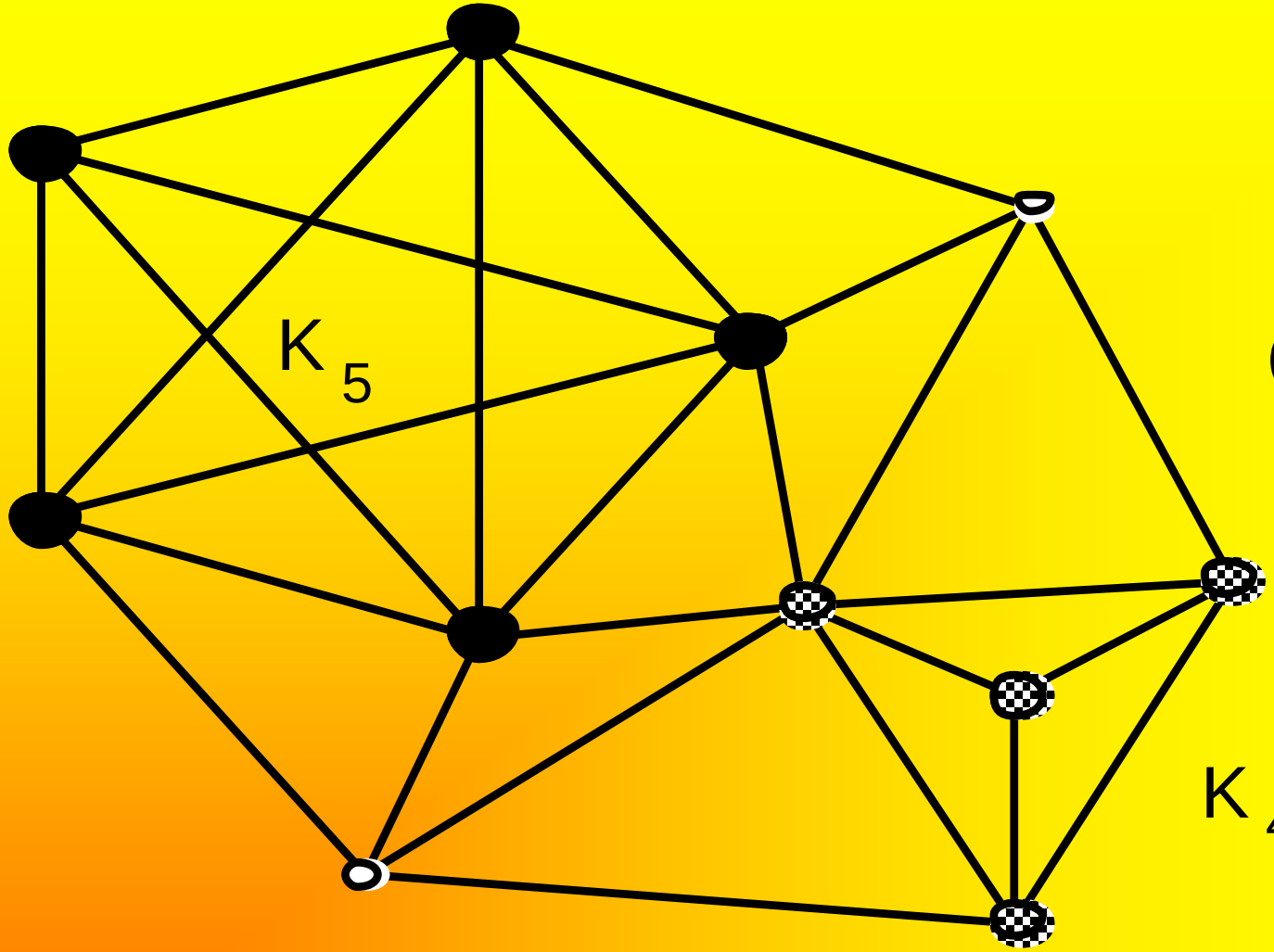
VÄLIVERKKO

Line graph



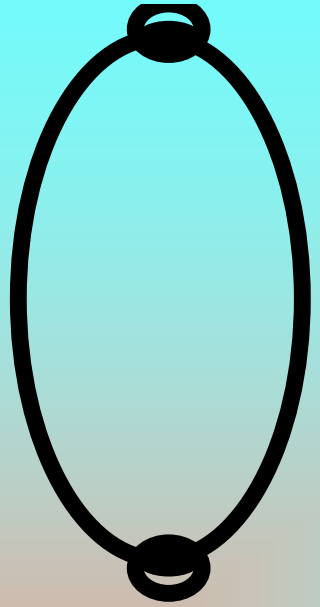
TÄYDELLINEN VERKKO

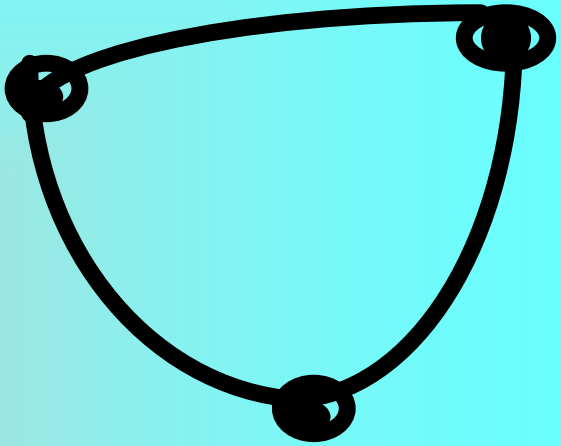
Kuratowski



KLIKKI

Klikkiluku d(G)





LUUPPI
RINNAKKAISET VÄLIT
KEHÄ

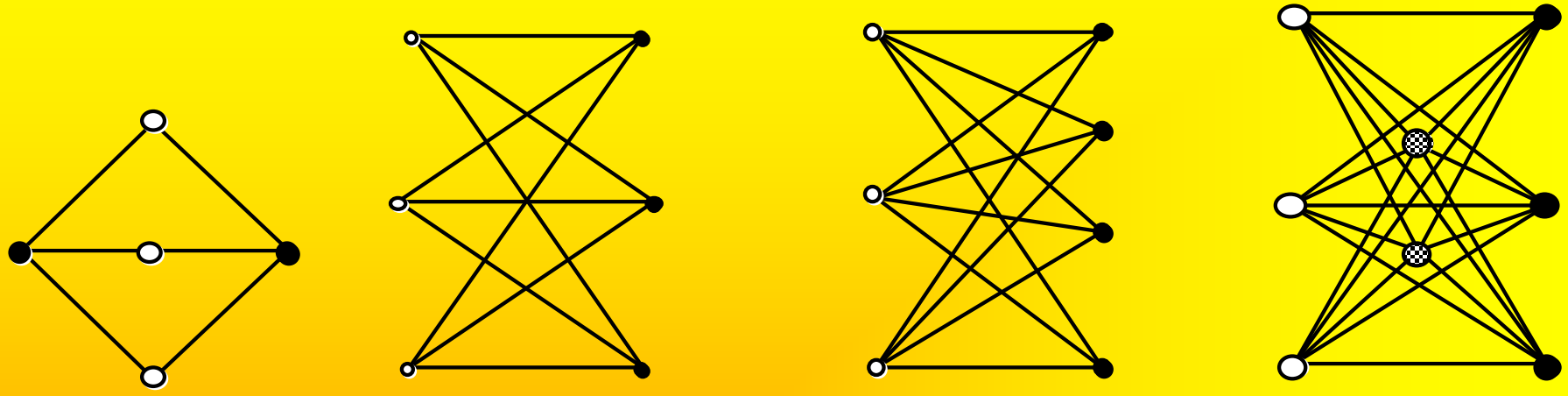
Closed path

AVOIN	SULJETTU	
T I E *W A L K*	SULJETTU TIE *CLOSED WALK*	koostuu väleistä ja solmuista tie alkaa aina solmusta ja päättyy solmuun
K E T J U *T R A I L*	S I L M U K K A *CLOSED TRAIL* *(CIRCUIT)*	jokainen väli esiintyy vain kerran
P O L K U *P A T H*	SULJETTU POLKU *CLOSED PATH*	jokainen solmu (ja väli) esiintyy vain kerran

erityisesti suljettu polku on:

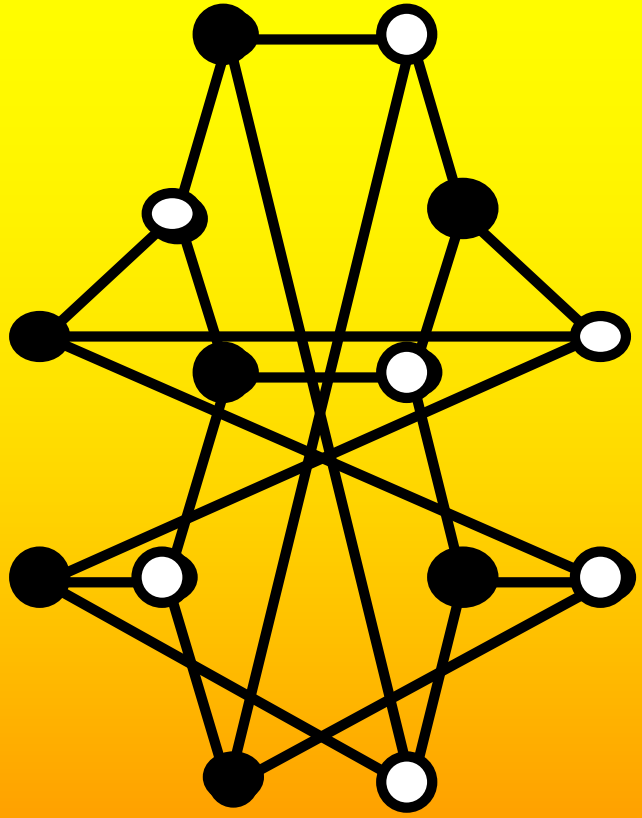
kun polun pituus on:

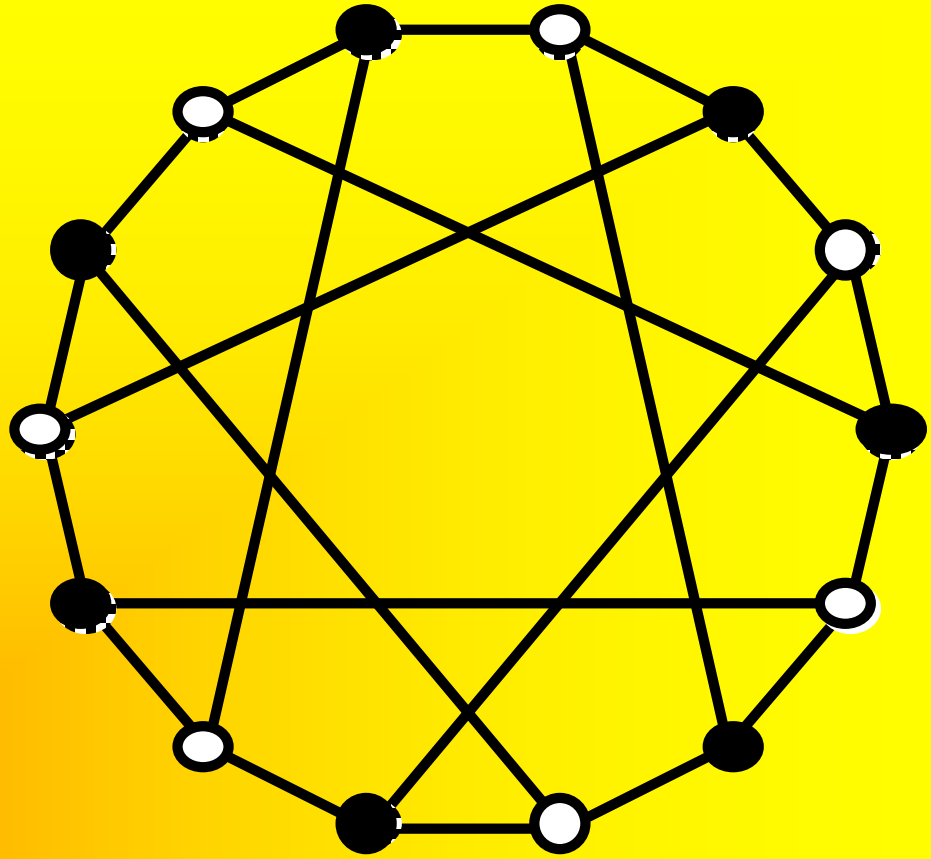
KEHÄ *CYCLE*	•3
LUUPPI (LENKKI) *LOOP*	1



k - JAKOINEN VERKKO

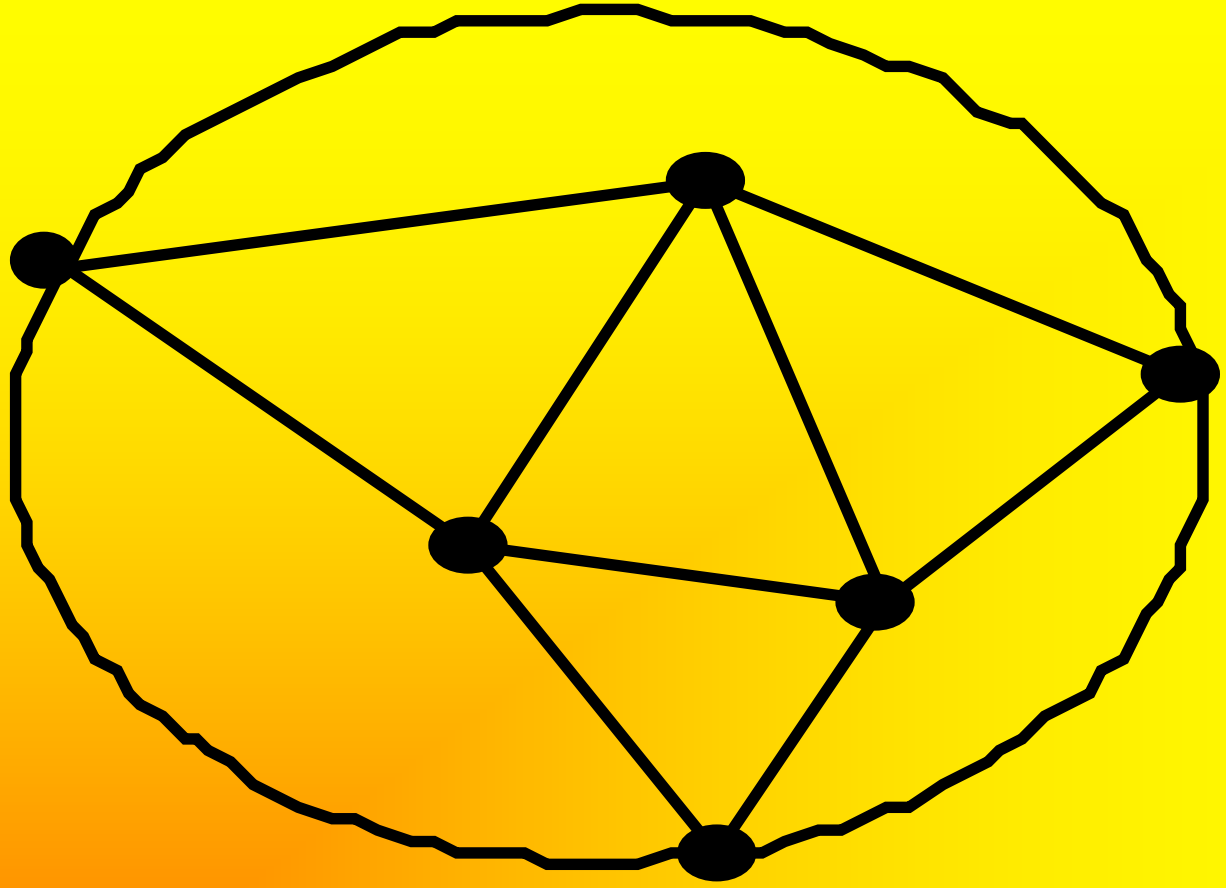
Bipartite & $K_{3,2,3}$





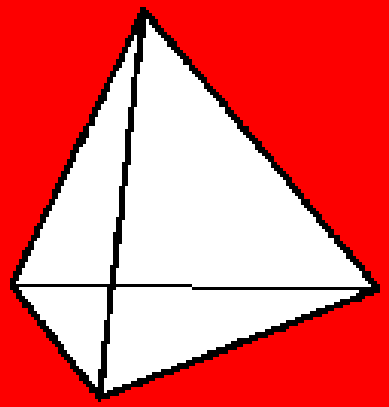
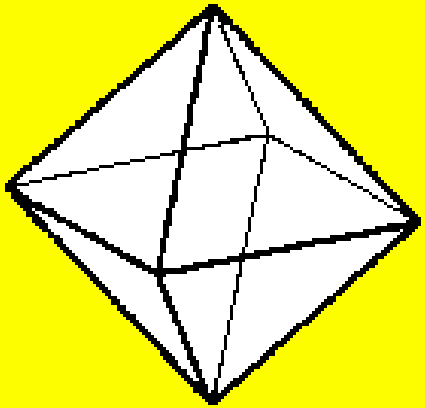
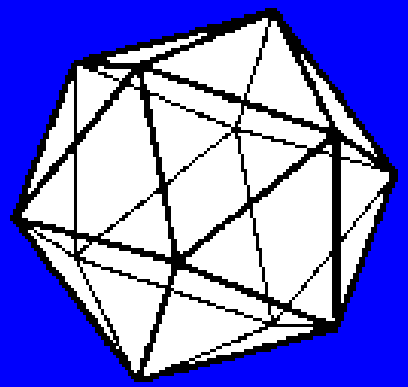
TRIVALENTTI

Bipartite & 2-yhtenäinen Heawood

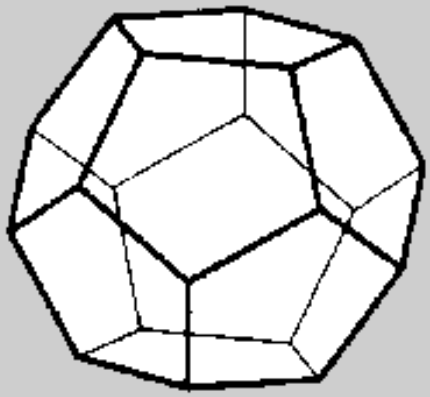


k - yhtenäinen verkko

$$V - E + F = 2$$

Fire	Air	Water
 Tetrahedron	 Octahedron	 Icosahedron
hot & dry	hot & wet	cold & wet

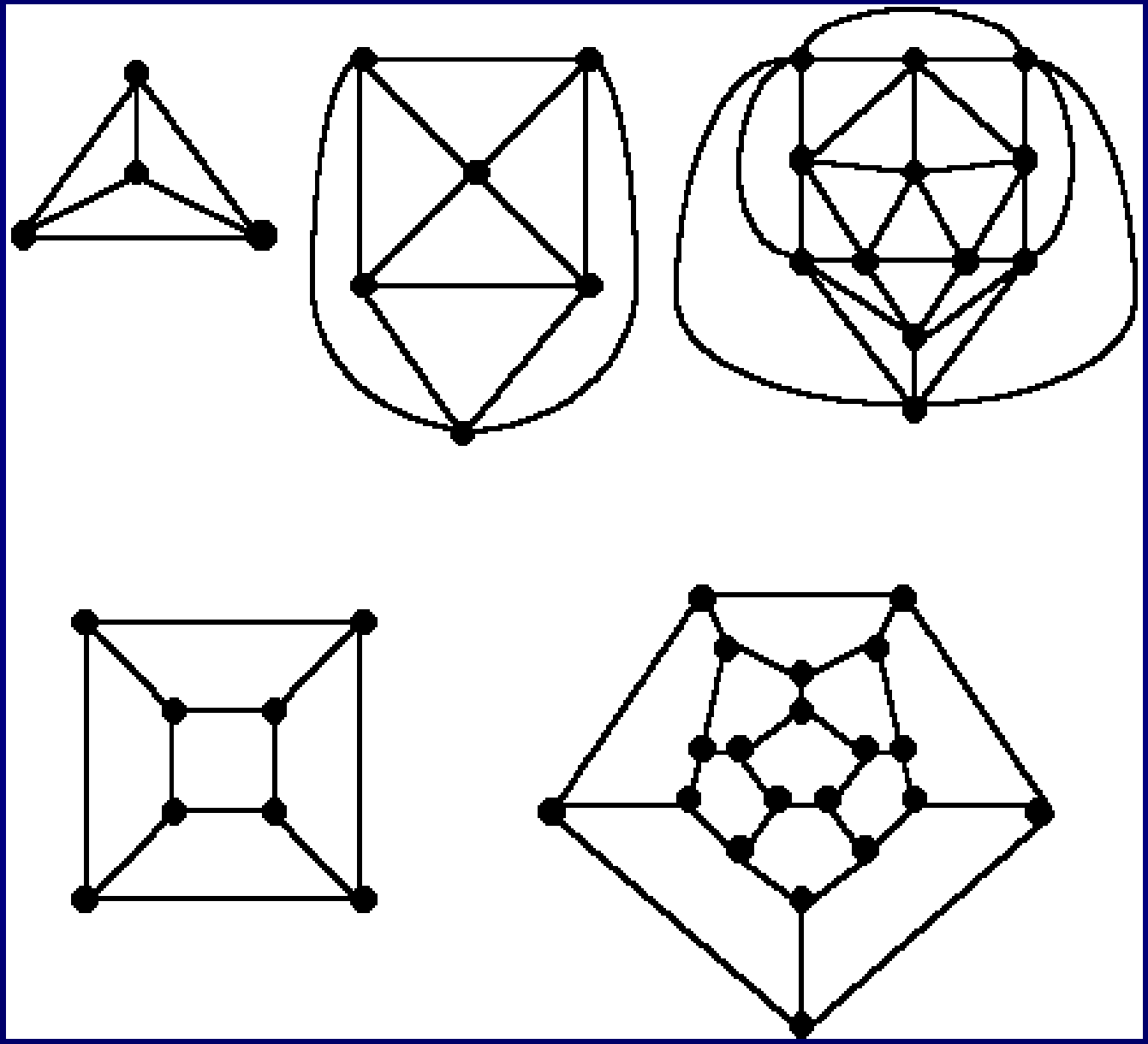
Empedocles	Earth
Plato: Platonic Solid with Squares	 Cube
Aristotle	cold & dry

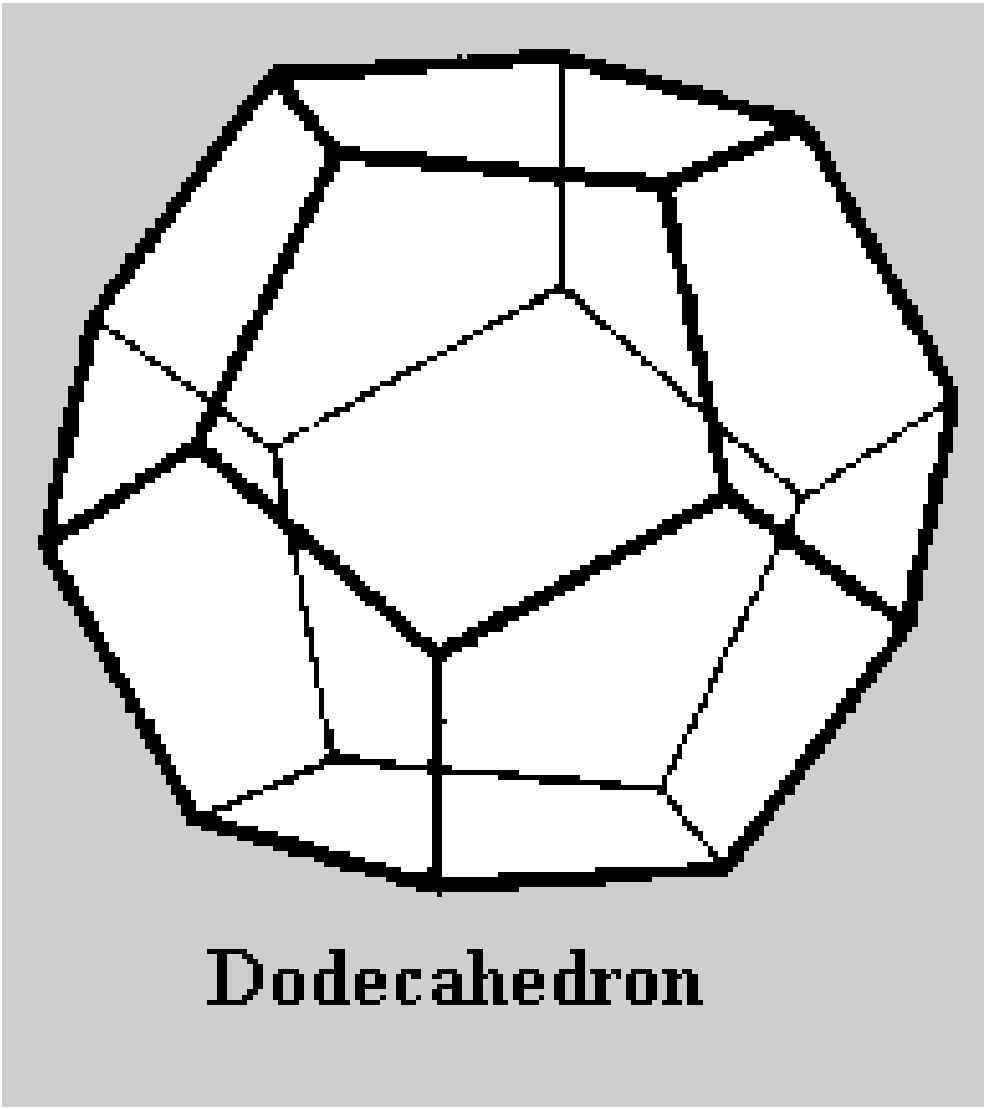


Dodecahedron

back

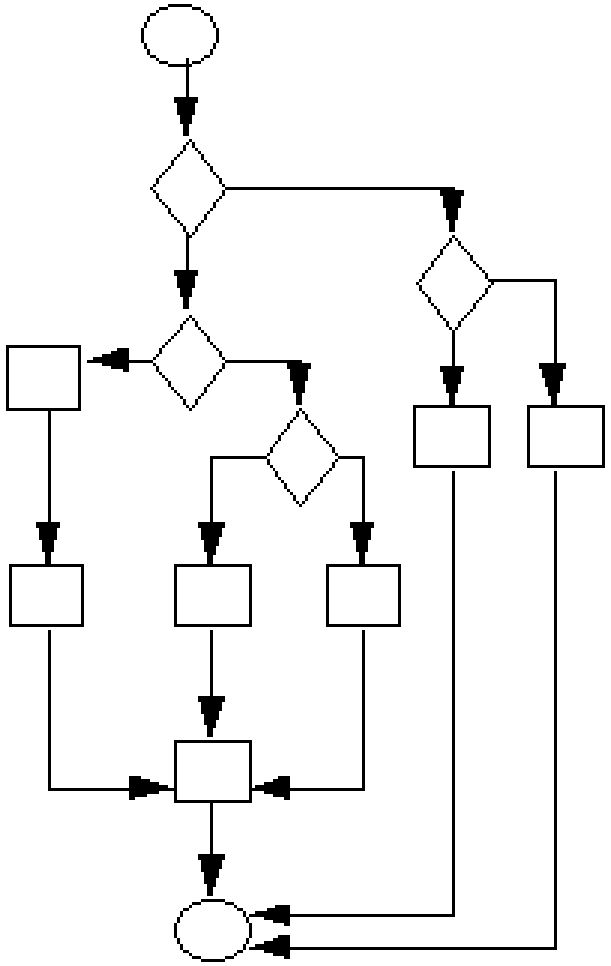
Netspace





Dodecahedron

McCabe and the cyclomatic complexity in programming



If G is the control flowgraph of program P and G has e edges (arcs) and n nodes

$$v(P) = e-n+2$$

v (P) is the number of linearly independent paths in G

here e = 16 n =13 v (P) = 5

More simply, if d is the number of decision nodes in G then

$$v(P) = d+1$$

Cyclomatic number, $v(G)$

For a graph G

$$v(G) = e - n + 2p$$

where

n is the number of vertices,

e is the number of edges,

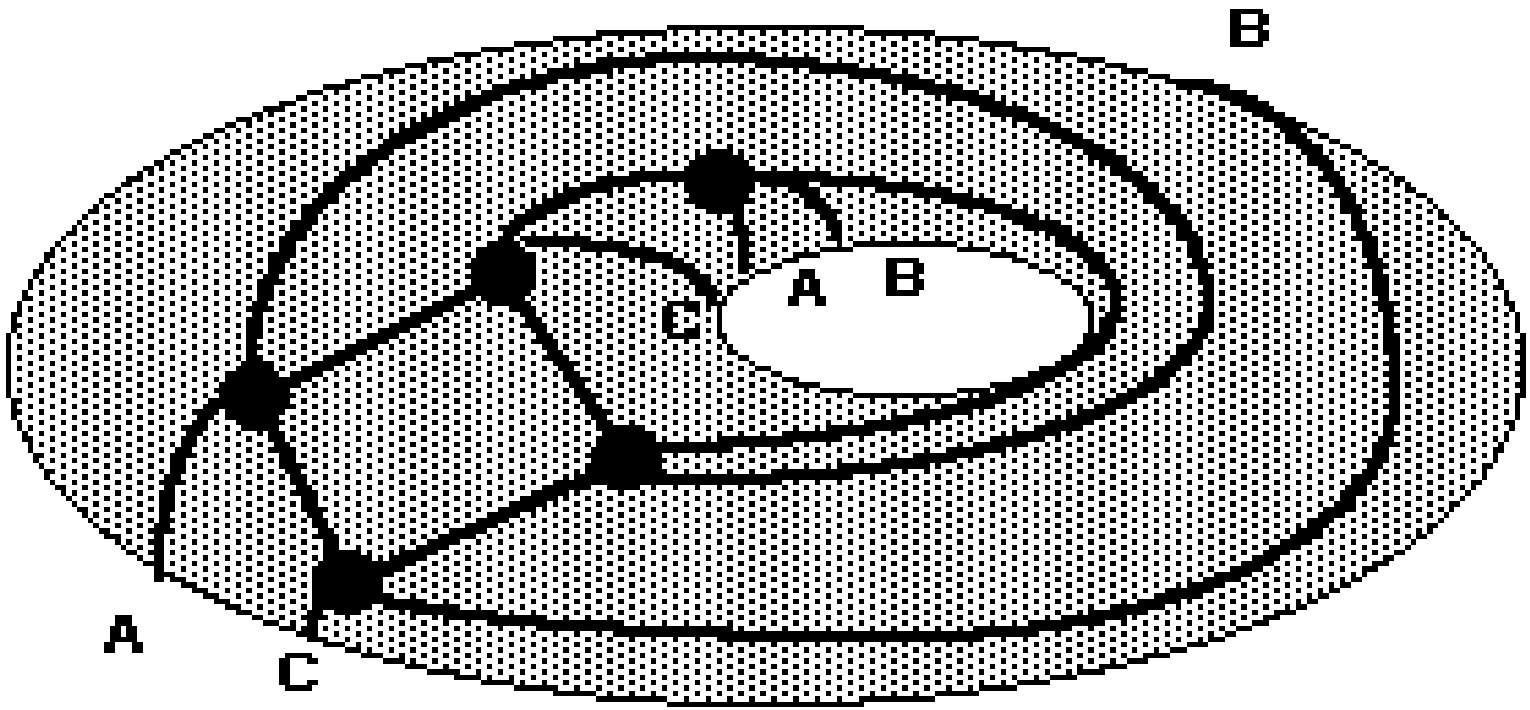
p is the number of strongly connected components.

back Netspace

$$E = 3V - 6$$

4CC

Netspace


B
A
B
C
A
C